

Hex Rules & Implementation

CMPUT 355: Games, Puzzles, and Algorithms

Lecture Outline

1. Logistics & Recap
2. Rules of Hex
3. Implementation Issues

Logistics

- **TA Office hours:** Every **Thursday** from **1:00pm-2:00pm** in **UCOMM 3-136**
 - Drop in basis; just show up and ask questions
 - Started yesterday (how'd it go?)
- **Practice quiz questions:** Released later today (**Jan 16**)
 - Answers released Tuesday (Jan 20)
- **Quiz 1:** Friday, **Jan 23**
 - In-class, full 50 minutes
 - No need to email if you have to miss it; up to 3 replaced by final exam automatically
 - Questions will be very similar to practice questions
 - (at least 3 will be *suspiciously* similar!)

Logistics: Code walkthroughs

Question: Which style of code walkthrough did you prefer?

```
go_helper.py [master]

156 def makemove(self, where, color):
157     assert (self.brd[where] == EMPTY), 'that point is not empty'
158
159     self.brd = change_string(self.brd, where, color)
160
161     cap = []
162     for j in self.nbr_offsets:
163         x = where + j
164         if self.brd[x] == opponent(color):
165             cap += self.captured(x, opponent(color))
166
167     if (len(cap)>0):
168         #print('removing captured group at', point_to_alphanum(wh
169         ere, self.C))
170         for j in cap:
171             self.brd = change_string(self.brd, j, EMPTY)
172         return cap, True # move ok so far
173
174     if self.captured(where, color):
175         print('whoops no liberty there: not allowed')
```

hexgo/stone_board.py

Track groups and liberties instead of searching after each move:

```
self.stones = [set(), set()] # start with empty board
self.blocks = {} # point (block name) -> stones in block
self.nbrs = {} # point (block name) -> neighbors
self.liberties = {} # point (block name) -> liberties
self.parents = {} # point -> parent in block
```

Use find and union to update tracking after each move:

def **add_stone**(self, color, point):

```
self.stones[color].add(point)
self.blocks[point].add(point)
```

for n in self.nbrs[point]:

```
if n in self.stones[color]: # same-color nbr
    self.merge_blocks(n, point)
```

```
if n in self.stones[Cell.opponent(color)]: # opponent nbr
    self.remove_liberties(n, point)
```

Union to update block membership:

def **merge_blocks**(self, p, q):

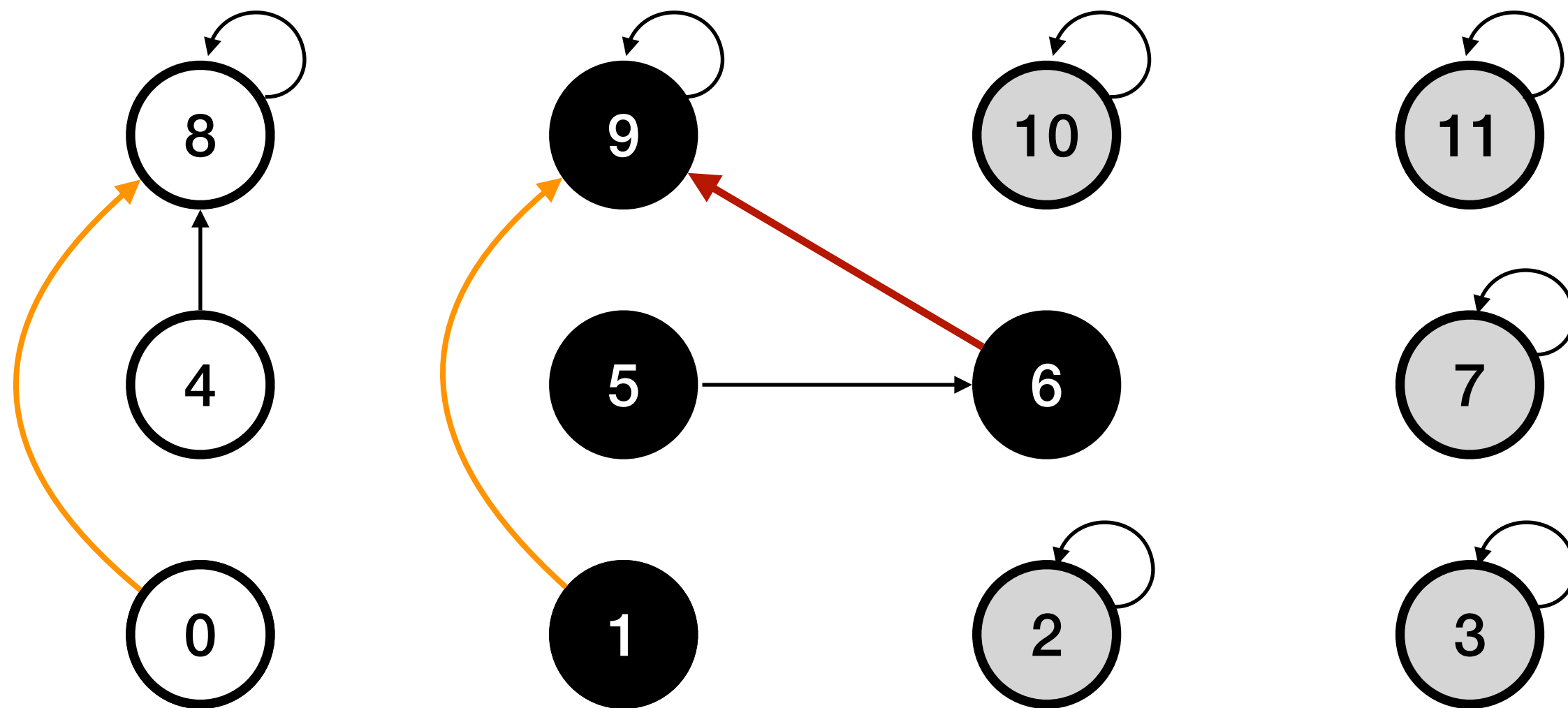
```
proot, qroot = UF.union(self.parents, p, q)
self.blocks[proot].update(self.blocks[qroot])
self.liberties[proot].update(self.liberties[qroot])
self.liberties[proot] -= self.blocks[proot]
```

Find to update block liberties:

def **remove_liberties**(self, p, q):

```
proot = UF.find(self.parents, p)
qroot = UF.find(self.parents, q)
self.liberties[proot] -= self.blocks[qroot]
self.liberties[qroot] -= self.blocks[proot]
```

Recap: Union-Find Operations in Go



After 5 is placed, we check neighbours 6, 9, 4, 1 for merge operations

1. Join 5's block (i.e., 5) to 6's block, so it points to 6's block's root (i.e., 6)
2. Next, join 5's block (i.e., 6) and 9's block (i.e., 9).
3. Without the union-by-rank optimization, that could mean *either* 6 points to 9 or 9 points to 6; I chose to show 6 points to 9.
4. But with the union-by-rank optimization, we would have to make 9 point to 6, because it has strictly lower rank than 6.

- Black stone on 6
- White stone on 8
- Black stone on 9
- White stone on 4
- Black stone on 5
- White stone on 0
- Black stone on 1
 - 0's block's liberties become empty

Recap: hexgo/stone_board.py

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Union to update block membership:

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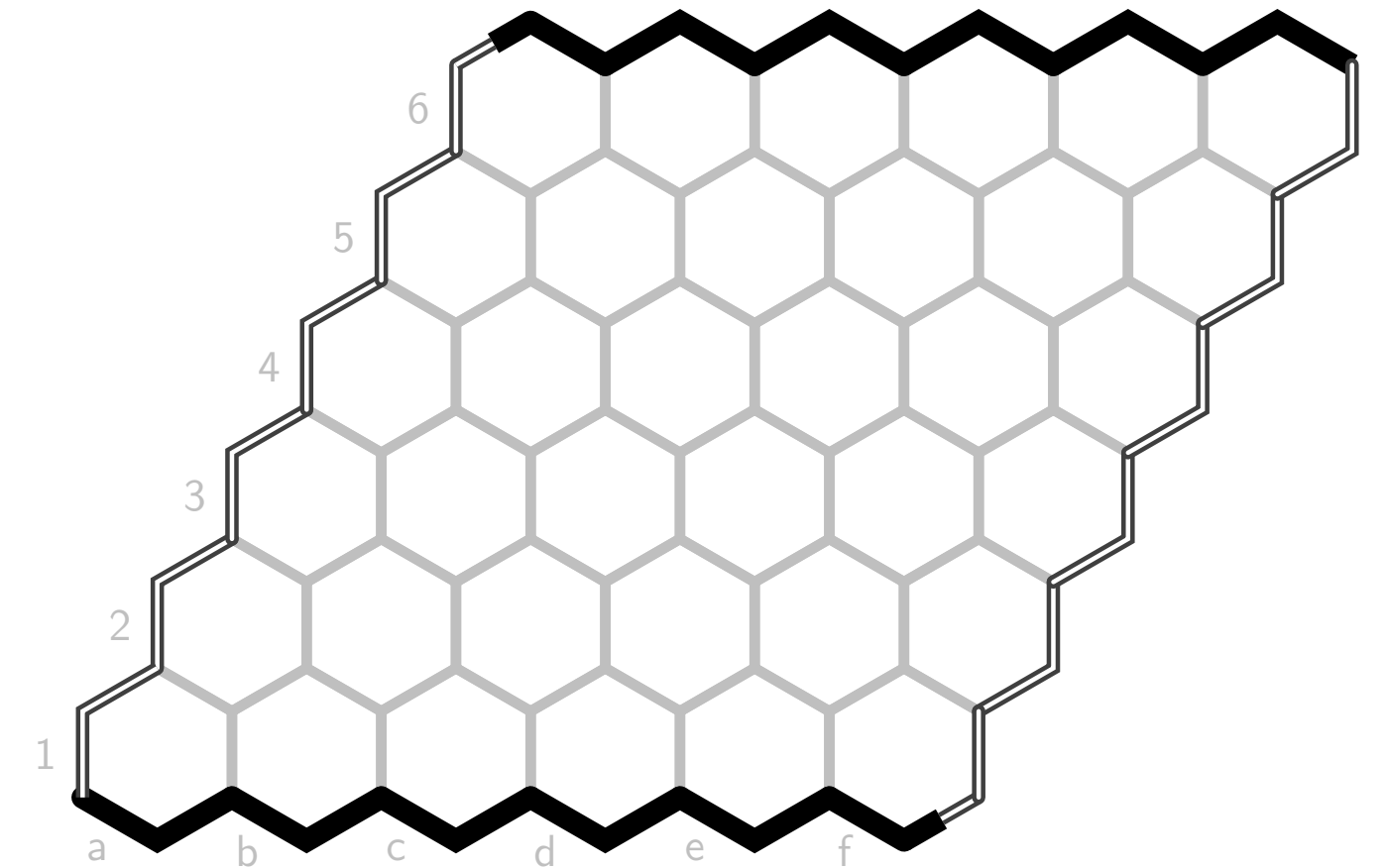
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def remove_liberties(self, p, q):
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```

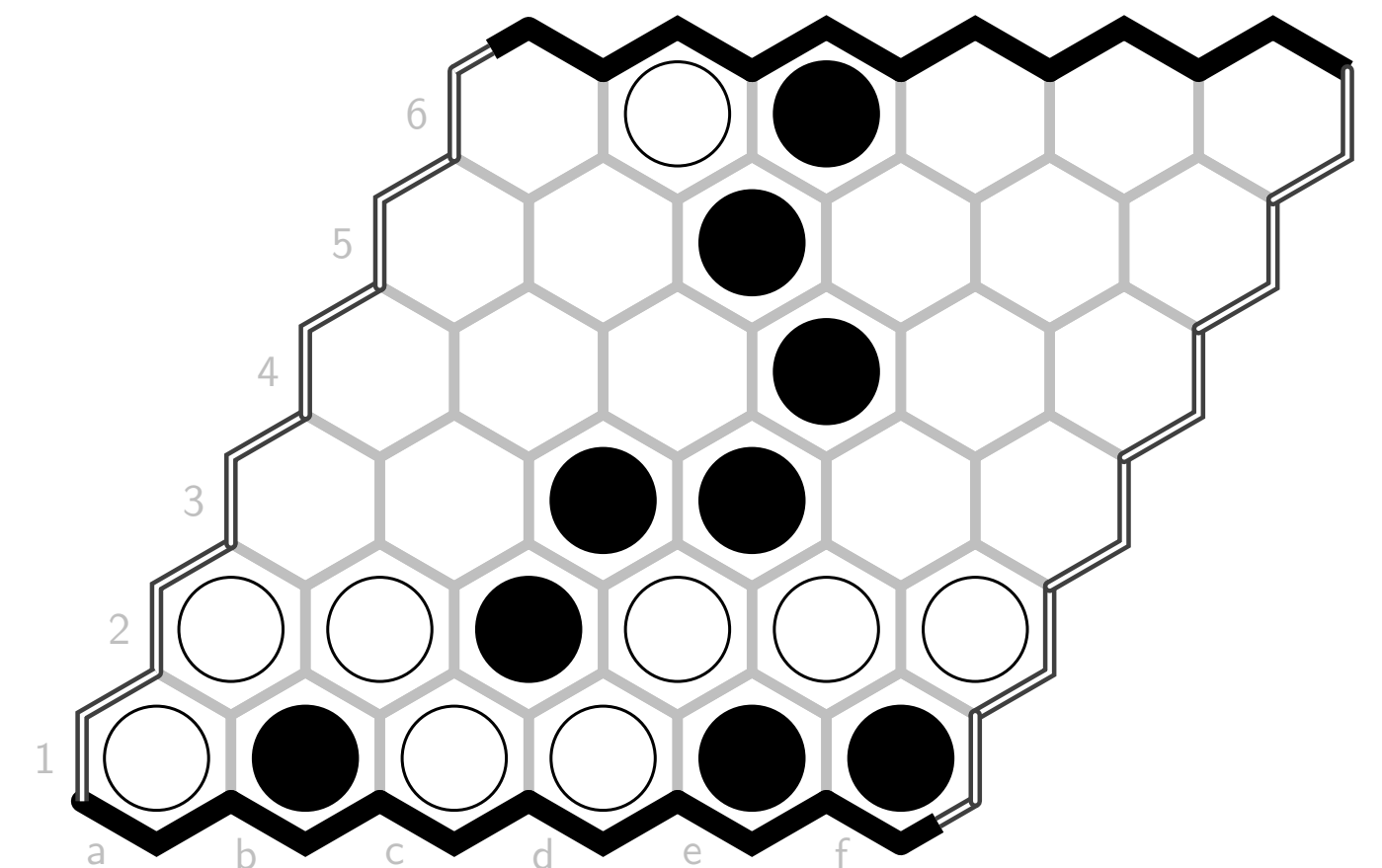

Hex Rules

- There are two players, Black and White
- Each player makes a **move** in alternation, starting with Black
 - A move is placing a stone in an Empty hexagonal cell
- Two cells are **adjacent** if they share a side
 - Each cell has 2-6 neighbours
- Two facing **borders** of the board are Black, the other two edges are White
 - Each bottom cell is adjacent to the bottom border, etc.
 - Each **corner** cell is thus adjacent to **two** borders
- The game ends when one player has joined the two edges of their own colour with a path of stones
 - The player who joins their edges **wins**

Empty Hex board

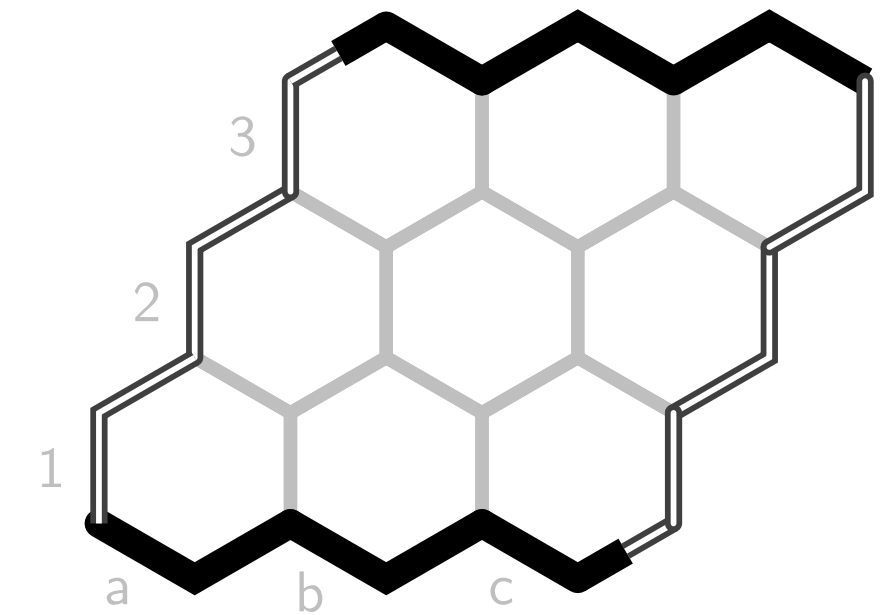
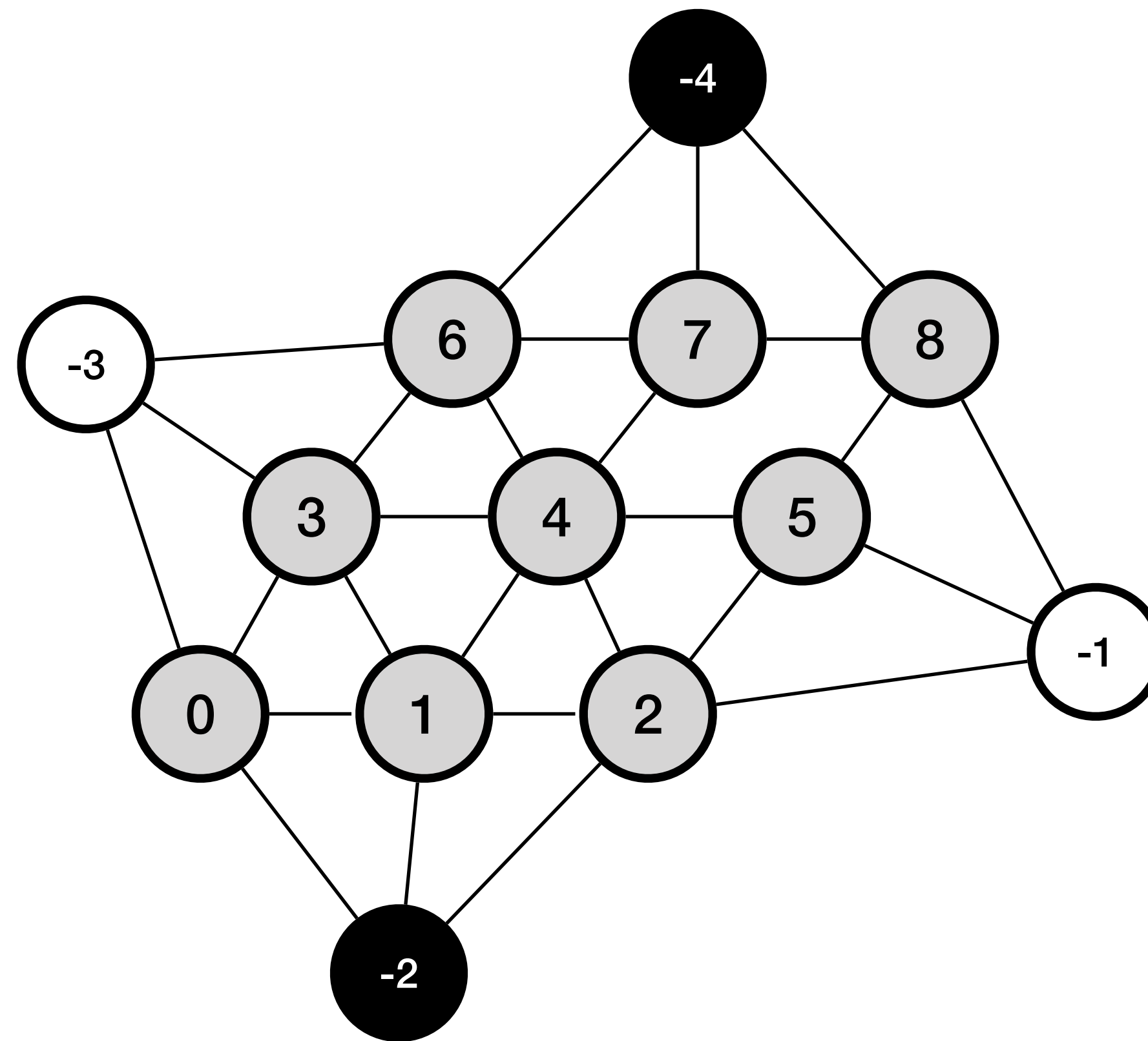


Winning position for White



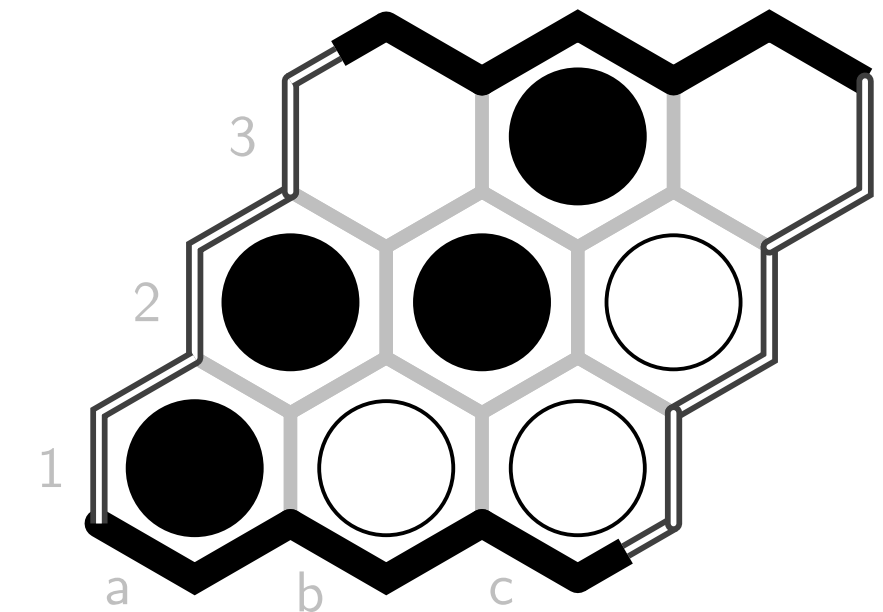
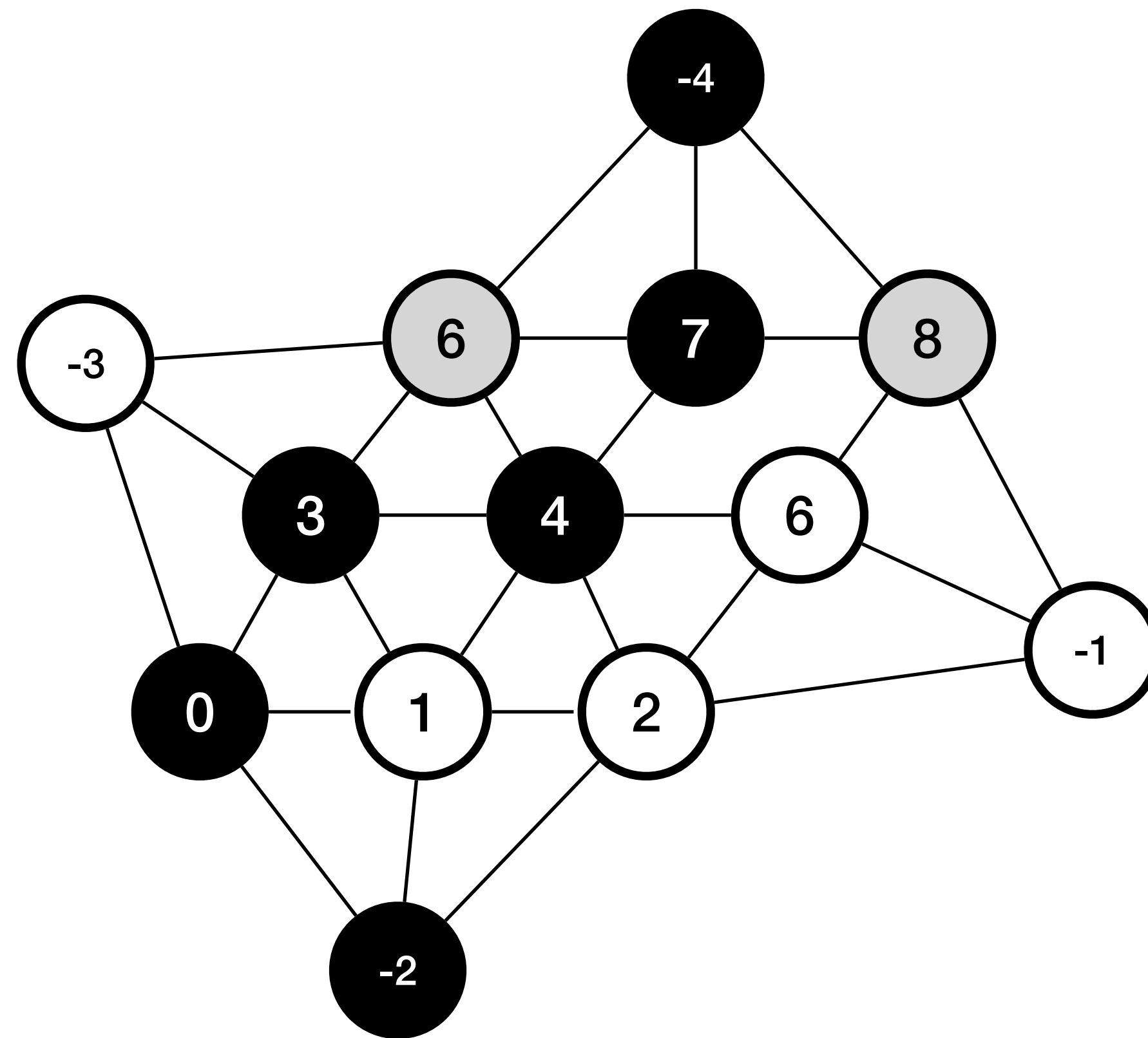
Graph Representation

- As with Go, can represent the board as a graph
 - One node per hex cell
 - Edge between each pair of adjacent cells
 - Plus four extra nodes for borders
- Each node can be coloured Black, White, or Empty
- **Question:** How to check for a winning condition on the graph?



Graph Representation

- As with Go, can represent the board as a graph
 - One node per hex cell
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Neighbours in hexgo/stone_board.py

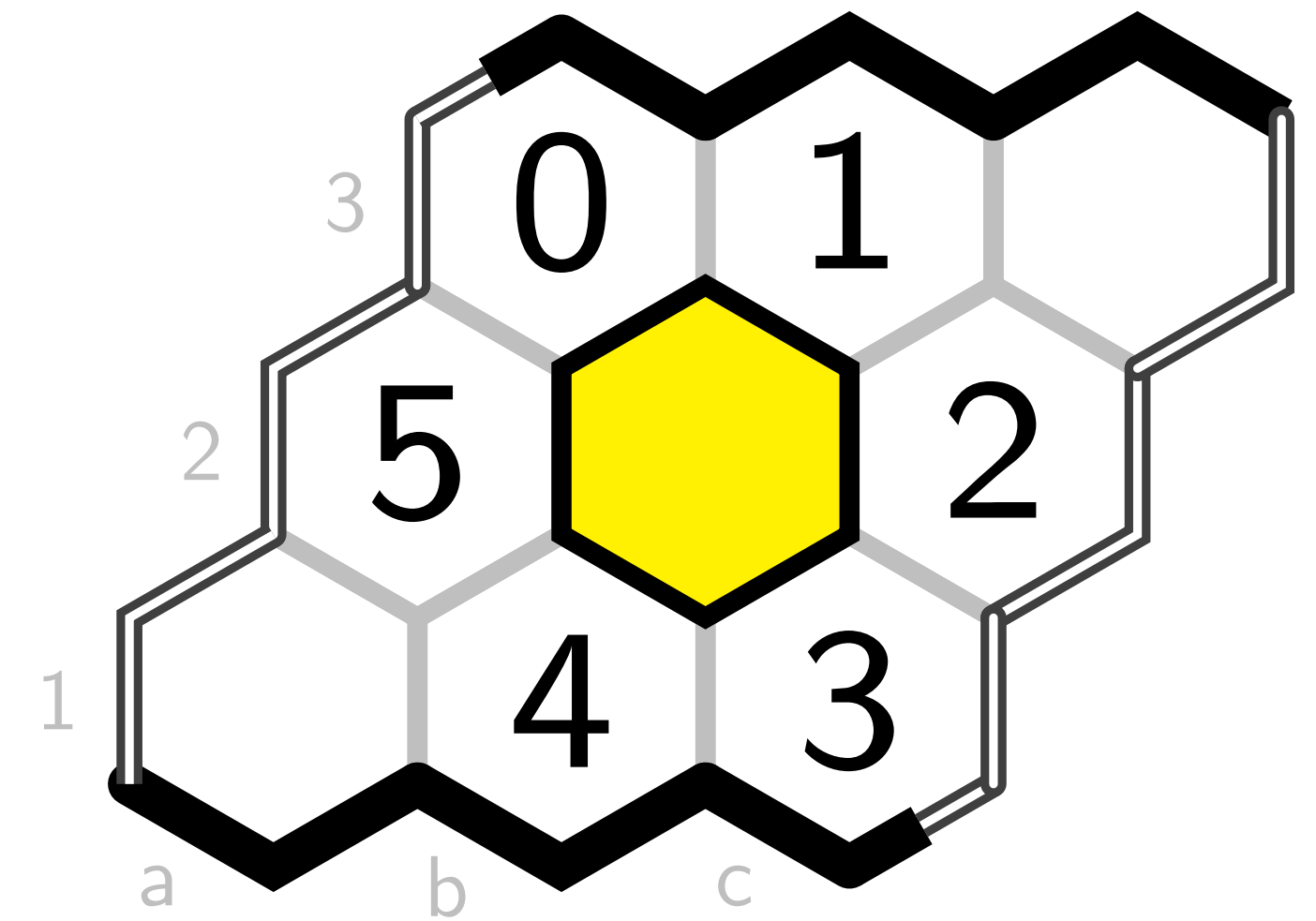
Instead of calculating a cell's neighbours every time we need them, we compute all of them **in advance** and then look them in `self.nbrs` whenever we need them. The following code is all in `Stone_board.__init__`:

Each cell has (up to) six neighbours:

```
if gt: # hex game
    self.top, self.rgt, self.btm, self.lft = -4, -3, -2, -1
    self.border = range(self.top, 0) # -4, -3, -2, -1
    self.p_range = range(self.top, self.n) # -4, ..., rows*cols-1
    self.nbr_offset = ((-1,0),(-1,1),(0,1),(1,0),(1,-1),(0,-1))
    # 0 1
    # 5 . 2
    # 4 3
```

Create an empty set to store each cell and border's neighbours:

```
for point in self.p_range:
    self.nbrs[point] = set()
```



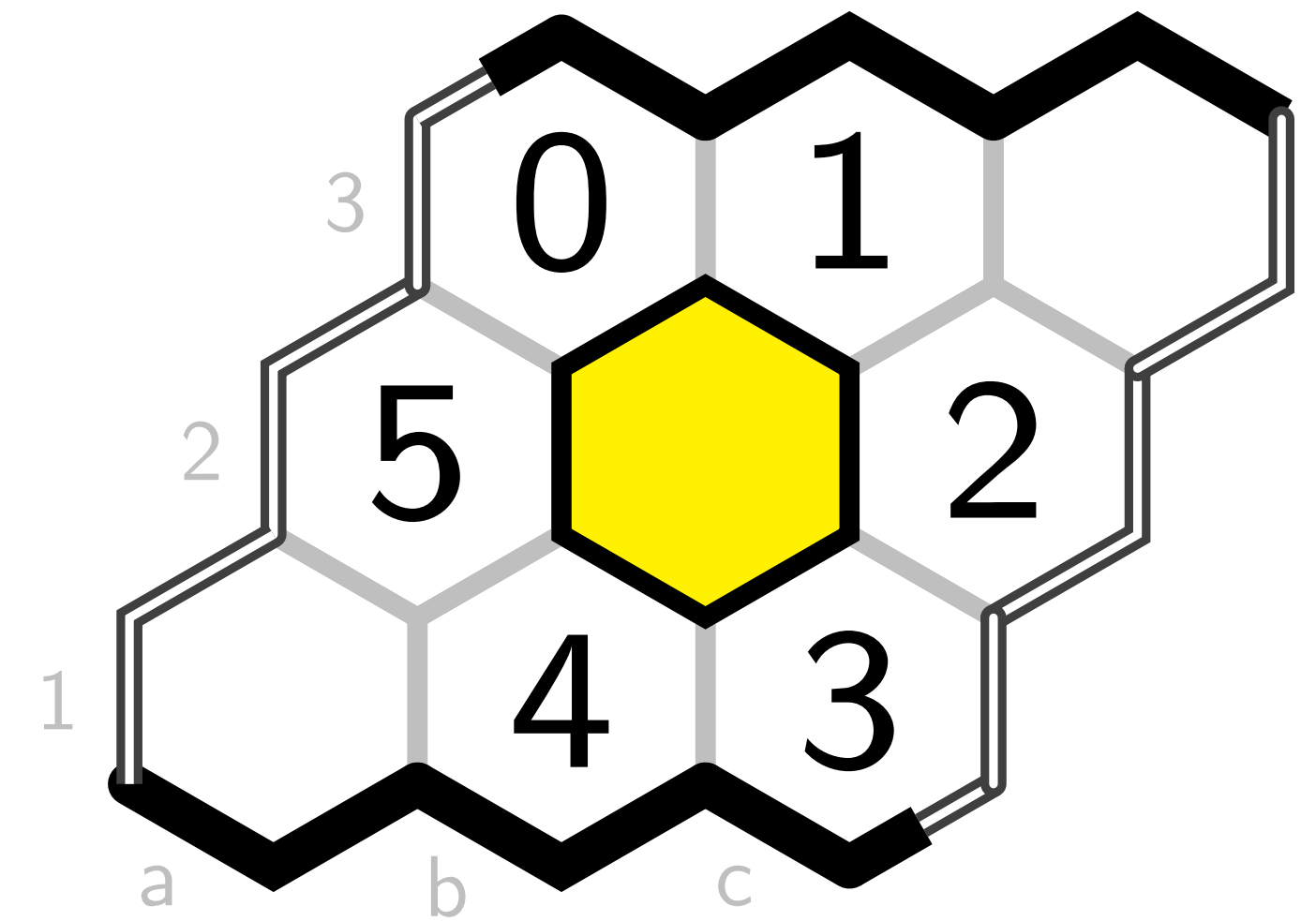
Initializing Neighbours

First record all the neighbouring **cells**:

```
r_range, c_range = range(self.r), range(self.c)
for r in range(self.r):
    for c in range(self.c):
        for (y,x) in self.nbr_offset:
            if r+y in r_range and c+x in c_range:
                self.nbrs[Pt.rc_point(r,c,self.c)].add(Pt.rc_point(r+y,c+x,self.c))
```

Then the neighbouring **borders**:

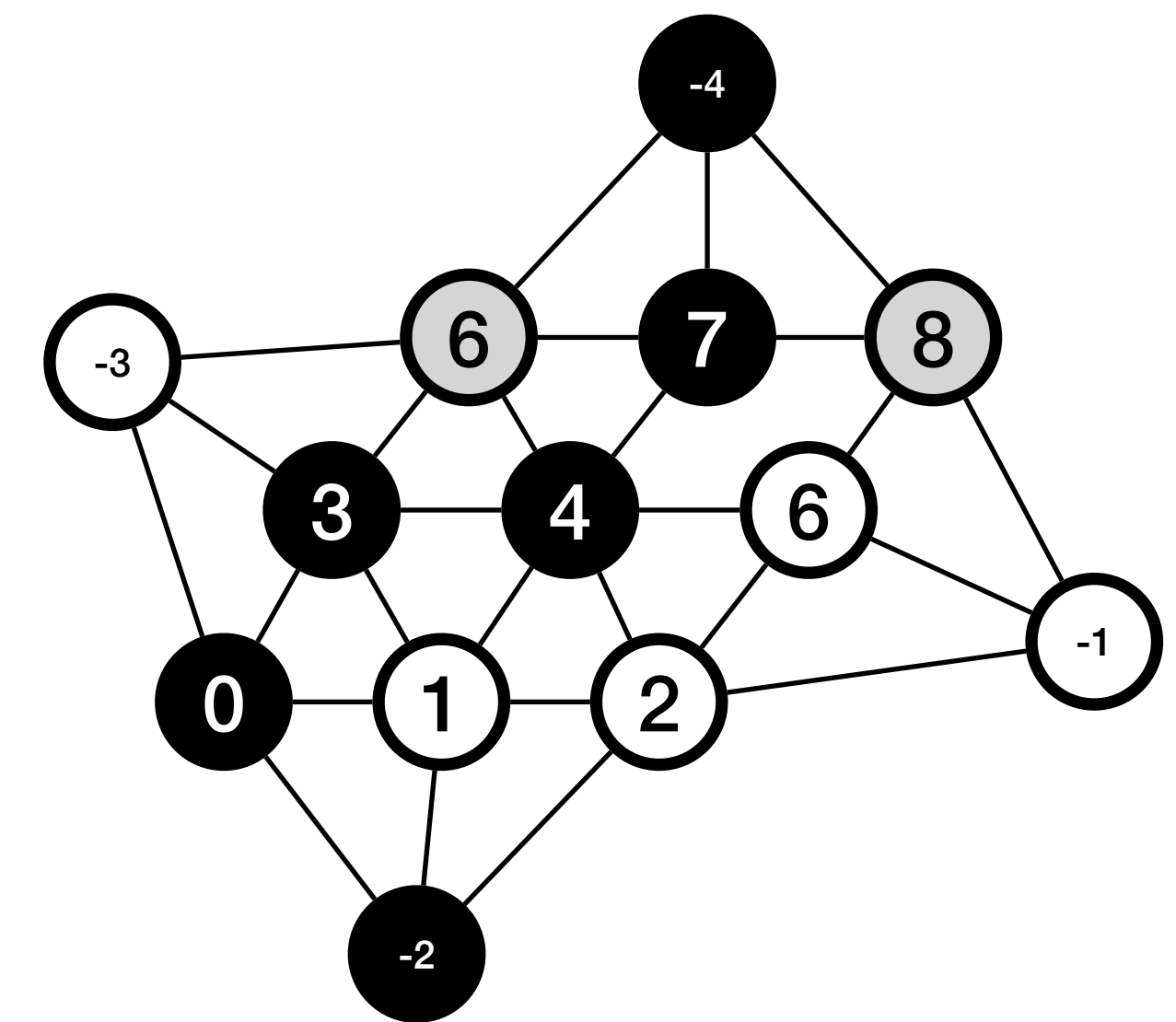
```
for j in range(self.c):
    self.nbrs[self.top].add(j)
    self.nbrs[j].add(self.top)
    self.nbrs[self.btm].add(self.n-j-1)
    self.nbrs[self.n-j-1].add(self.btm)
for k in range(self.r):
    self.nbrs[self.lft].add(k*self.c)
    self.nbrs[k*self.c].add(self.lft)
    self.nbrs[self.rgt].add(k*self.c+self.c-1)
    self.nbrs[k*self.c+self.c-1].add(self.rgt)
```



Paths

- A path in hex is just a group of adjacent cells with the same colour
 - It's a **block**!
- Can use *exactly* the same union-find approach to track paths as in Go:

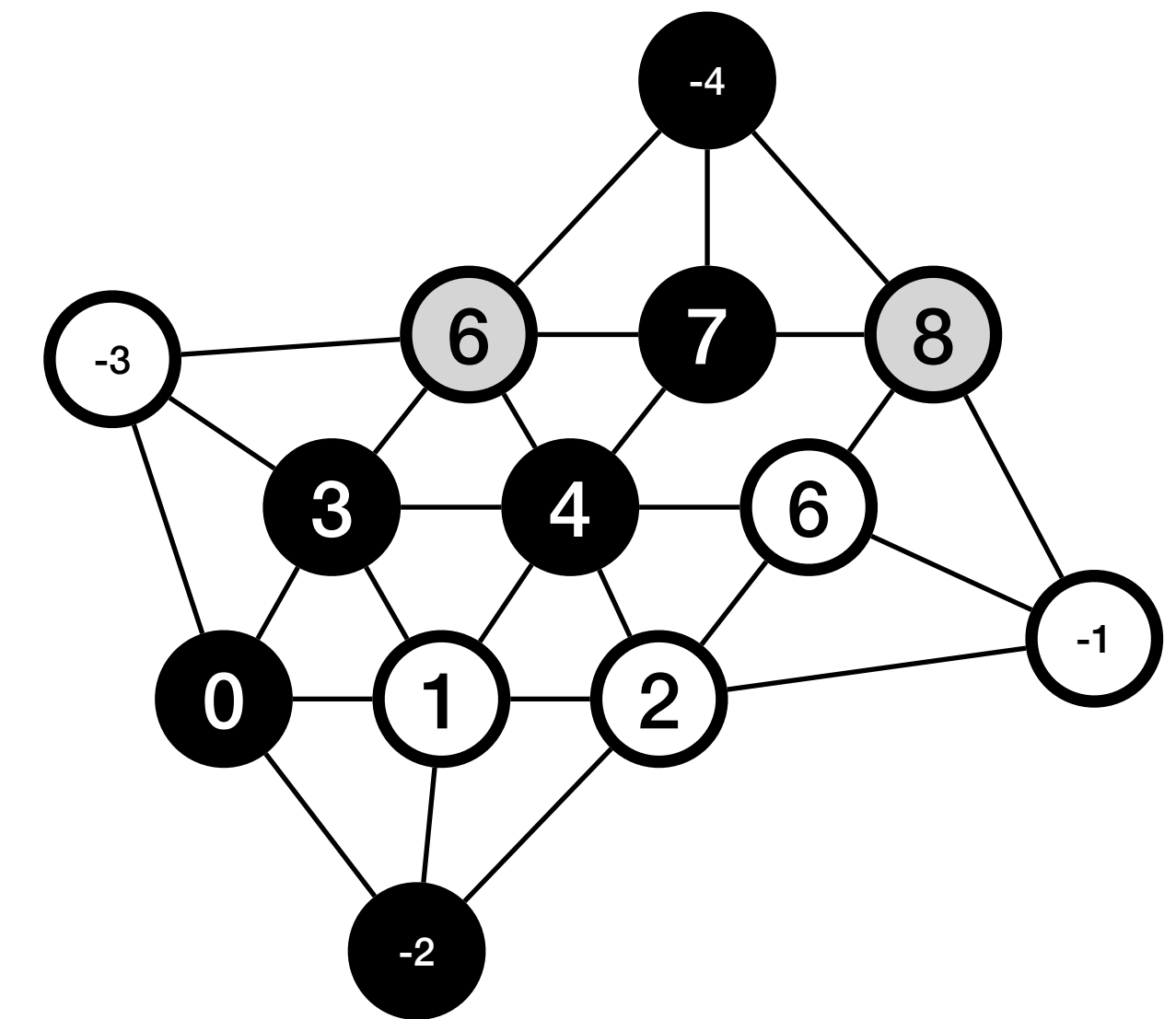
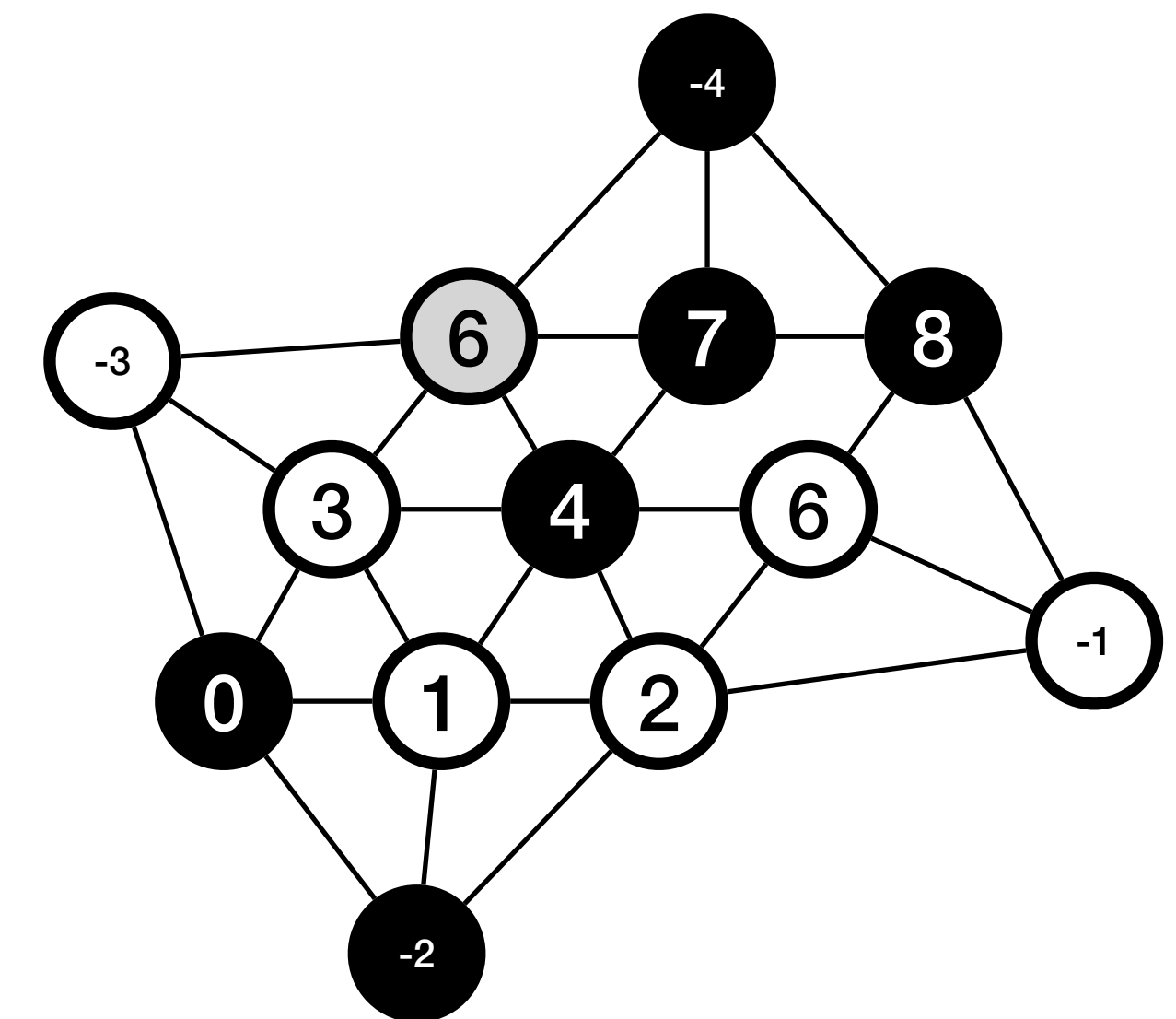
```
def add_stone(self, color, point):  
    self.stones[color].add(point)  
    self.blocks[point].add(point)  
  
    for n in self.nbrs[point]:  
        if n in self.stones[color]: # same-color nbr  
            self.merge_blocks(n, point)  
        if n in self.stones[Cell.opponent(color)]: # opponent nbr  
            self.remove_liberties(n, point)
```



End of Game

- **Question:** How can we detect the end of the game using our union-find approach?
 - **Black win:** There is a path from the top to the bottom, so there must be a block that contains both the top and bottom borders
 - **White win:** Must be a block containing both the left and right borders
- So just check if top and bottom (or left and right) are in the same block:

```
def hex_win(self, cell_color):  
    if self.game_type != Game.hex_game:  
        return False  
    if cell_color == Cell.b:  
        return UF.in_same_block(self.parents, self.top, self.btm)  
    return UF.in_same_block(self.parents, self.lft, self.rgt)
```



Summary

- Rules of Hex:
 - Players alternately place stones in hexagonal cells
 - First player to connect their two borders with a path of stone of their colour wins
- Implementing Hex in `hexgo/stone_board.py`:
 - Same graph representation as Go:
Each point gets an index on a (sort of) rectangular grid
 - Borders represented by special "cells" that have negative indices
 - Paths represented by blocks tracked with Union-Find datastructure (exactly the same code as Go!)
 - Win condition: Both of a player's borders are in the same block