

Social Choice

CMPUT 654: Modelling Human Strategic Behaviour

S&LB §9.1-9.4

Lecture Outline

1. Recap
2. Risk Aversion
3. Aggregating Preferences
4. Voting Paradoxes
5. Arrow's Theorem

Recap: Bayesian Games

- **Epistemic types** are a profile of signals that **parameterize the utility functions** of each agent
 - Possibly correlated
 - Each agent observes only their own type
- Three notions of **expected utility**:
 - **ex-ante**: before observing type
 - **ex-interim**: after observing own type
 - **ex-post**: full type profile is known
- Solution concepts:
 - **Bayes-Nash equilibrium**: equilibrium of induced normal form of **ex-ante utilities**
 - **Ex-post equilibrium**: Agents are best-responding at **every type profile** (not just in expectation)

Familiar-Looking Question

- Consider the lottery [.9999: \$1, .0001: \$10,000,000]
Expected value: \$1000.9999
- **Question:** Can a rational expected utility maximizer with **monotonic value for money** turn down a ticket to this lottery that costs \$100?

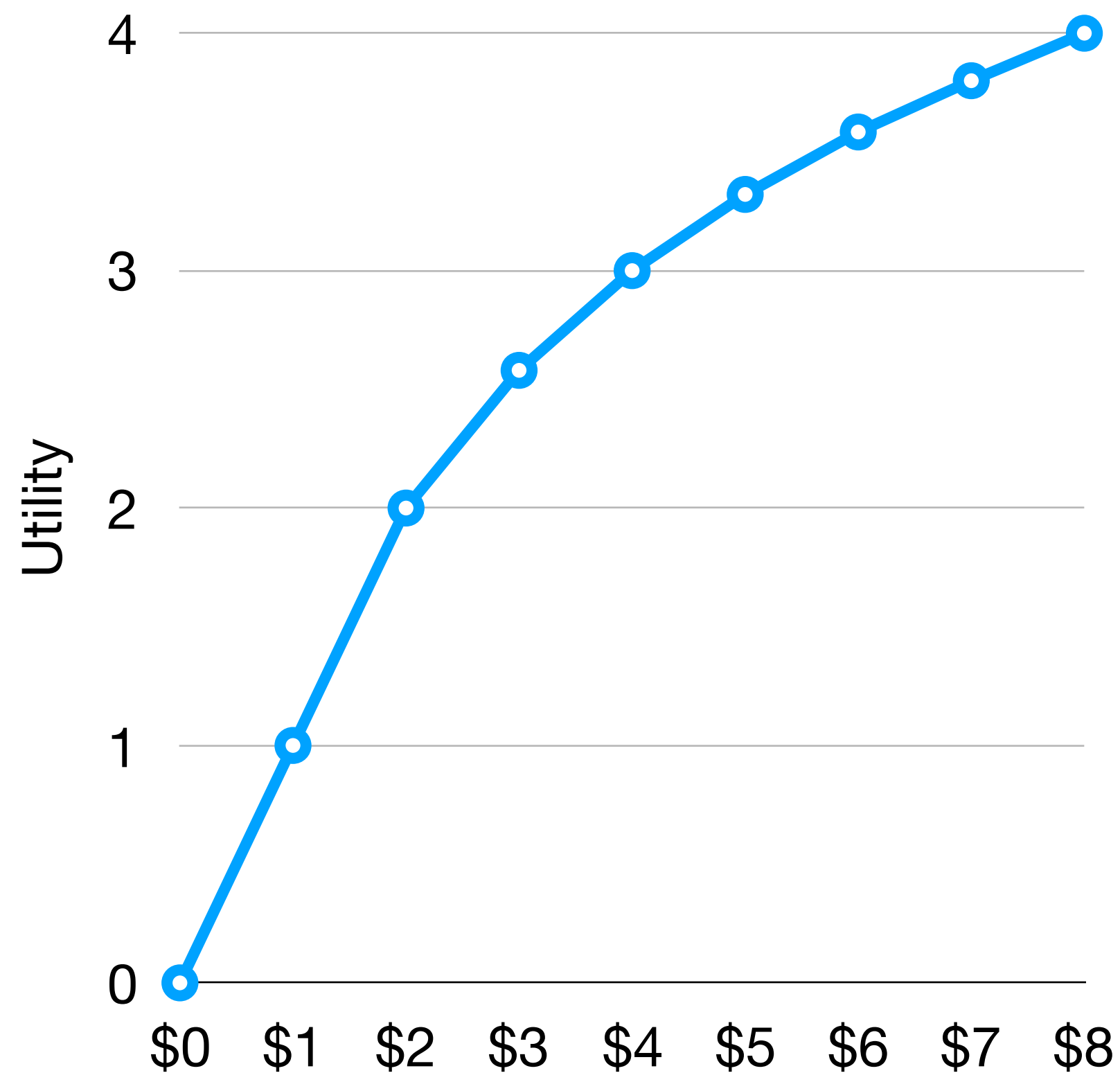
Risk Attitudes

- Rational agents are required to maximize their **expected utility**
- This is distinct from the **expected value** of some outcome (e.g., money)
- Rational agents can have a utility for money that is:
 - **Risk-averse:** Value of a lottery is **lower** than expected value
 - **Risk-neutral:** Value of a lottery is **exactly** expected value
 - **Risk-seeking:** Value of a lottery is **higher** than expected value

Risk Aversion Example

- Expectations are taken over **outcomes**:
 - $\mathbb{E}[u(o)] \neq u(\mathbb{E}[o])$
- Example: $\ell = [.5 : \$0, .5 : \$8]$ has expected **value** of \$4

$$\begin{aligned} u(\ell) &= \mathbb{E}[u(o)] = .5u(\$0) + .5u(\$8) \\ &= 2 = .5(0) + .5(4) \\ &\neq 3 = u(\$4) = u(\mathbb{E}[o]) \end{aligned}$$



Risk Attitudes and Utility for Money

- **Concave** utility for money \Rightarrow **risk aversion**
- **Linear** utility for money \Rightarrow **risk neutral**
- **Convex** utility for money \Rightarrow **risk seeking**

Aggregating Preferences

- Suppose we have a collection of agents, each with **individual preferences** over some **outcomes**
- How should we choose the outcome?
 - Ignore strategic reporting issues: Either the center already knows everyone's preferences, or the agents don't lie
- **More formally:** Can we construct a **social choice function** that maps the profile of preference orderings to an outcome?
- **More generally:** Can we construct a **social welfare function** that maps the profile of preference orderings to an **aggregated preference ordering**?

Formal Model

Definition: A **social choice function** is a function $C : L^n \rightarrow O$, where

- $N = \{1, 2, \dots, n\}$ is a set of **agents**
- O is a finite set of **outcomes**
- L is the set of non-strict total orderings over O .

Definition: A **social welfare function** is a function $C : L^n \rightarrow L$, where N , O , and L are as above.

Notation:

We will denote **i 's preference order** as $\succeq_i \in L$, and a **profile** of preference orders as $[\succeq] \in L^n$.

Non-ranking Voting Schemes

Voters need not submit a full preference ordering:

1. **Plurality voting:** Everyone votes for favourite outcome, choose the outcome with the most votes
2. **Cumulative voting:** Everyone is given k votes to distribute among candidates as they like; choose the outcome with the most votes
3. **Approval voting:** Each agent casts a single vote for each of the outcomes that are "acceptable"; choose the outcome with the most votes.

Ranking Voting Schemes

Every agent expresses their **full preference ordering**:

1. **Plurality with elimination**

- Everyone votes for favourite outcome
- Outcome with least votes is eliminated
- Repeat until one outcome remains

2. **Borda**

- Everyone assigns scores to outcome: Most-preferred gets $n-1$, next-most-preferred gets $n-2$, etc. Least-preferred outcome gets 0.
- Outcome with highest sum of scores is chosen

3. **Pairwise Elimination**

- Define a schedule over the order in which pairs of outcomes will be compared
- For each pair, everyone chooses their favourite; least-preferred is eliminated
- Continue to next pair of non-eliminated outcomes until only one outcome remains

Condorcet Condition

Definition:

An outcome $o \in O$ is a **Condorcet winner** if $\forall o' \in O$,
 $|\{i \in N : o \succ_i o'\}| > |\{i \in N : o' \succ_i o\}|$.

Definition:

A social choice function satisfies the **Condorcet condition** if it always selects a Condorcet winner when one exists.

- If there's one outcome that would win a pairwise vote against every other possible outcome, then perhaps we want our social choice rule to pick it
- Unfortunately, such an outcome does not always exist
 - There can be cycles where A would beat B, B would beat C, C would beat A

Paradox: Condorcet Winner

499 agents: $a \succ b \succ c$

3 agents: $b \succ c \succ a$

498 agents: $c \succ b \succ a$

- **Question:** Who is the **Condorcet winner**?
- **Question:** Who wins a **plurality election**?
- **Question:** Who wins under **plurality with elimination**?

Paradox: Sensitivity to Losing Candidate

35 agents: $a \succ c \succ b$

33 agents: $b \succ a \succ c$

32 agents: $c \succ b \succ a$

- **Question:** Who wins under **plurality**?
- **Question:** Who wins under **Borda**?
- **Question:** Now **drop c**. Who wins under **plurality**?
- **Question:** After dropping c, who wins under **Borda**?

Paradox: Sensitivity to Agenda Setter

35 agents: $a \succ c \succ b$

33 agents: $b \succ a \succ c$

32 agents: $c \succ b \succ a$

- **Question:** Who wins under **pairwise elimination** with order **a, b, c** ?
- **Question:** Who wins with ordering **a, c, b** ?
- **Question:** Who wins with ordering **b, c, a** ?

Paradox: Sensitivity to Agenda Setter

1 agent: $b \succ d \succ c \succ a$

1 agent: $a \succ b \succ d \succ c$

1 agent: $c \succ a \succ b \succ d$

- **Question:** Who wins with ordering a, b, c, d ?
- **Question:** What is **wrong** with that?

Arrow's Theorem

These problems are not a coincidence; they affect **every possible** voting scheme.

Notation:

- For this section we switch to **strict total orderings** L
- Preference ordering selected by social welfare function W is $>_W$.

Pareto Efficiency

Definition:

W is **Pareto efficient** if for any $o_1, o_2 \in O$,

$$(\forall i \in N : o_1 \succ o_2) \implies (o_1 \succ_W o_2)$$

- If **everyone agrees** that o_1 is better than o_2 , then the aggregated preference order should reflect that

Independence of Irrelevant Alternatives

Definition:

W is independent of irrelevant alternatives if, for any $o_1, o_2 \in O$ and any two preference profiles $[>']$, $[>''] \in L$,

$$(\forall i \in N : o_1 \succ'_i o_2 \iff o_1 \succ''_i o_2) \implies (o_1 \succ_{W[>']} o_2 \iff o_1 \succ_{W[>'']} o_2)$$

- If every agent has the same ordering between two particular outcomes in two different preference profiles, then the social welfare function on those two profiles must order those two outcomes the same way
- The ordering between two outcomes should only depend on the agents' orderings between those outcomes, not on where any other outcomes are in the agents' orderings

Non-Dictatorship

Definition:

W does not have a **dictator** if

$$\neg i \in N : \forall [\succ] \in L^n : \forall o_1, o_2 \in O : (o_1 \succ_i o_2) \implies (o_1 \succ_W o_2)$$

- No single agent determines the social ordering

Arrow's Theorem

Theorem: (Arrow, 1951)

If $|O| > 2$, any social welfare function that is Pareto efficient and independent of irrelevant alternatives is dictatorial.

- Unfortunately, restricting to social choice functions instead of full social welfare functions doesn't help.

Theorem: (Muller-Satterthwaite, 1977)

If $|O| > 2$, any social choice function that is weakly Pareto efficient and monotonic is dictatorial.

Summary

- Social choice is the study of **aggregating** the true preferences of a group of agents
 - **Social choice function**: Chooses a **single outcome** based on preference profile
 - **Social welfare function**: Chooses a **full preference order** over outcomes based on preference profile
- Well-known voting rules all lead to **unfair or undesirable** outcomes
 - **Arrow's Theorem**: This is unavoidable
 - **Muller-Satterthwaite Theorem**: ... even restricting to social choice functions