

Imperfect Information Extensive Form Games

CMPUT 654: Modelling Human Strategic Behaviour

S&LB §5.2-5.2.2

Lecture Outline

1. Recap
2. Imperfect Information Games
3. Behavioural vs. Mixed Strategies
4. Perfect vs. Imperfect Recall
5. Computational Issues



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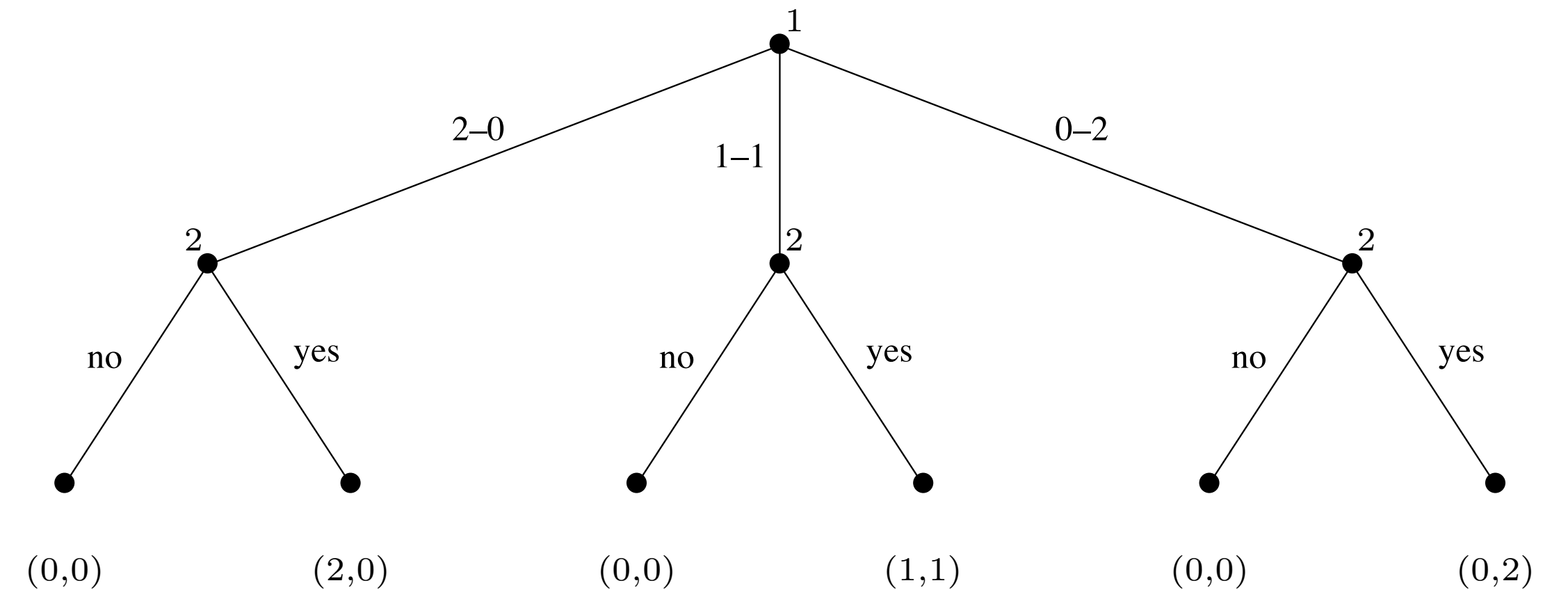
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Recap: Perfect Information Extensive Form Game

Definition:

A **finite perfect-information game in extensive form** is a tuple $G = (N, A, H, Z, \chi, \rho, \sigma, u)$, where

- N is a set of n **players**,
- A is a single set of **actions**,
- H is a set of nonterminal **choice nodes**,
- Z is a set of **terminal nodes** (disjoint from H),
- $\chi : H \rightarrow 2^A$ is the **action function**,
- $\rho : H \rightarrow N$ is the **player function**,
- $\sigma : H \times A \rightarrow H \cup Z$ is the **successor function**.
- $u = (u_1, u_2, \dots, u_n)$ is a **utility function** for each player $u_i : Z \rightarrow \mathbb{R}$.



Recap: Pure Strategies

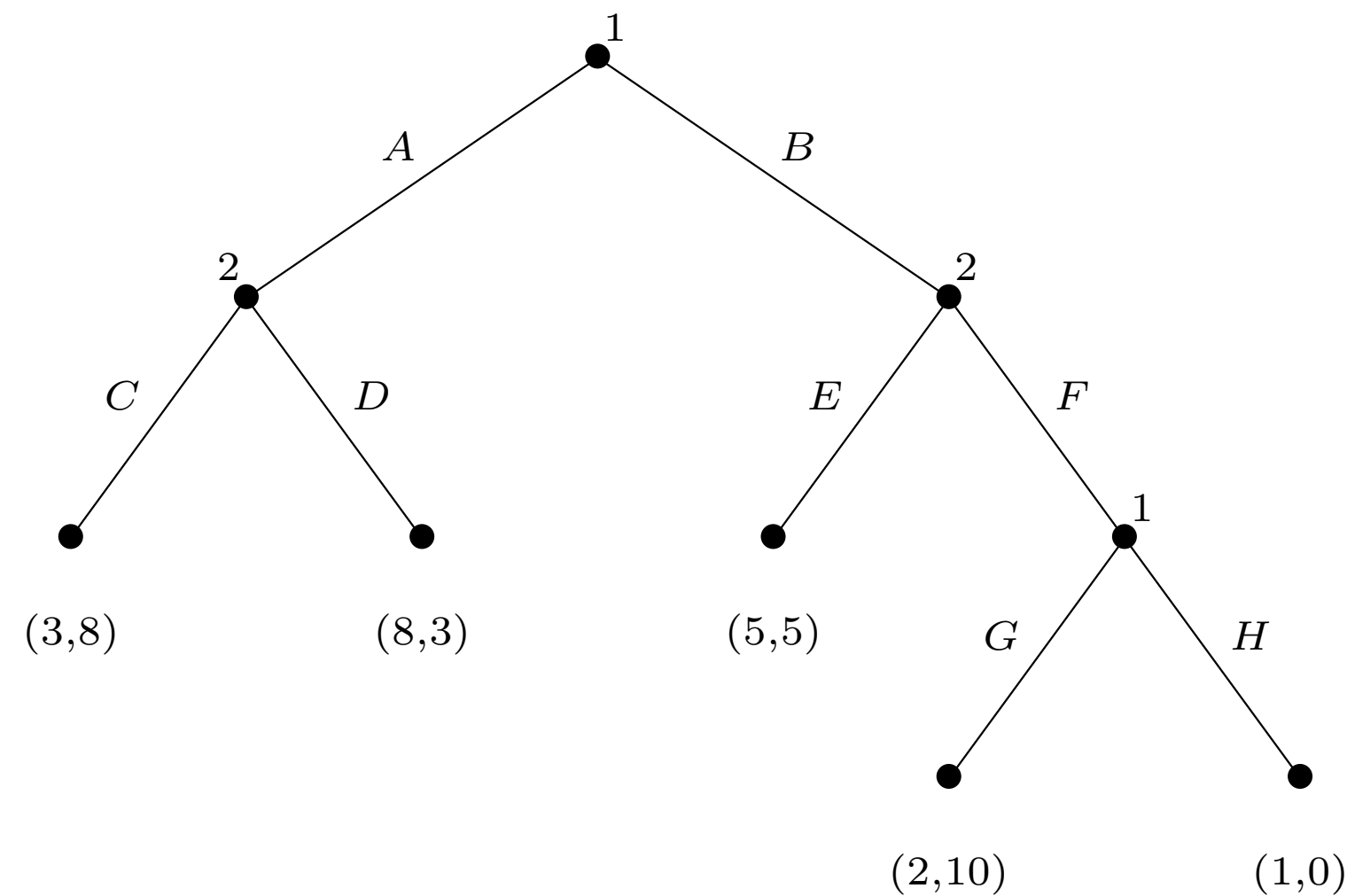
Definition:

Let $G = (N, A, H, Z, \chi, \rho, \sigma, u)$ be a perfect information game in extensive form. Then the **pure strategies of player i** consist of the cross product of actions available to player i at each of their choice nodes, i.e.,

$$\prod_{h \in H | \rho(h) = i} \chi(h)$$

- A pure strategy associates an action with **each** choice node, even those that will **never be reached**

Recap: Induced Normal Form



	C,E	C,F	D,E	D,F
A,G	3,8	3,8	8,3	8,3
A,H	3,8	3,8	8,3	8,3
B,G	5,5	2,10	5,5	2,10
B,H	5,5	1,0	5,5	1,0

- Any pair of pure strategies uniquely identifies a **terminal node**, which identifies a **utility** for each agent
- We have now defined a set of **agents**, **pure strategies**, and **utility functions**
- Any extensive form game defines a corresponding **induced normal form game**

Recap: Backward Induction

- **Backward induction** is a straightforward algorithm that is guaranteed to compute a subgame perfect equilibrium
- **Idea:** Replace subgames lower in the tree with their equilibrium values

```
BACKWARDINDUCTION( $h$ ):  
  if  $h$  is terminal:  
    return  $u(h)$   
   $i := \rho(h)$   
   $U := -\infty$   
  for each  $h'$  in  $\chi(h)$ :  
     $V = \text{BACKWARDINDUCTION}(h')$   
    if  $V_i > U_i$ :  
       $U_i := V_i$   
  return  $U$ 
```

Imperfect Information, informally

- **Perfect information** games model **sequential** actions that are **observed by all players**
 - **Randomness** can be modelled by a special *Nature* player with constant utility
- But many games involve **hidden** actions
 - Cribbage, poker, Scrabble
 - Sometimes actions of the **players** are hidden, sometimes **Nature's** actions are hidden, sometimes both
- **Imperfect information extensive form games** are a model of games with sequential actions, some of which may be **hidden**

Imperfect Information Extensive Form Game

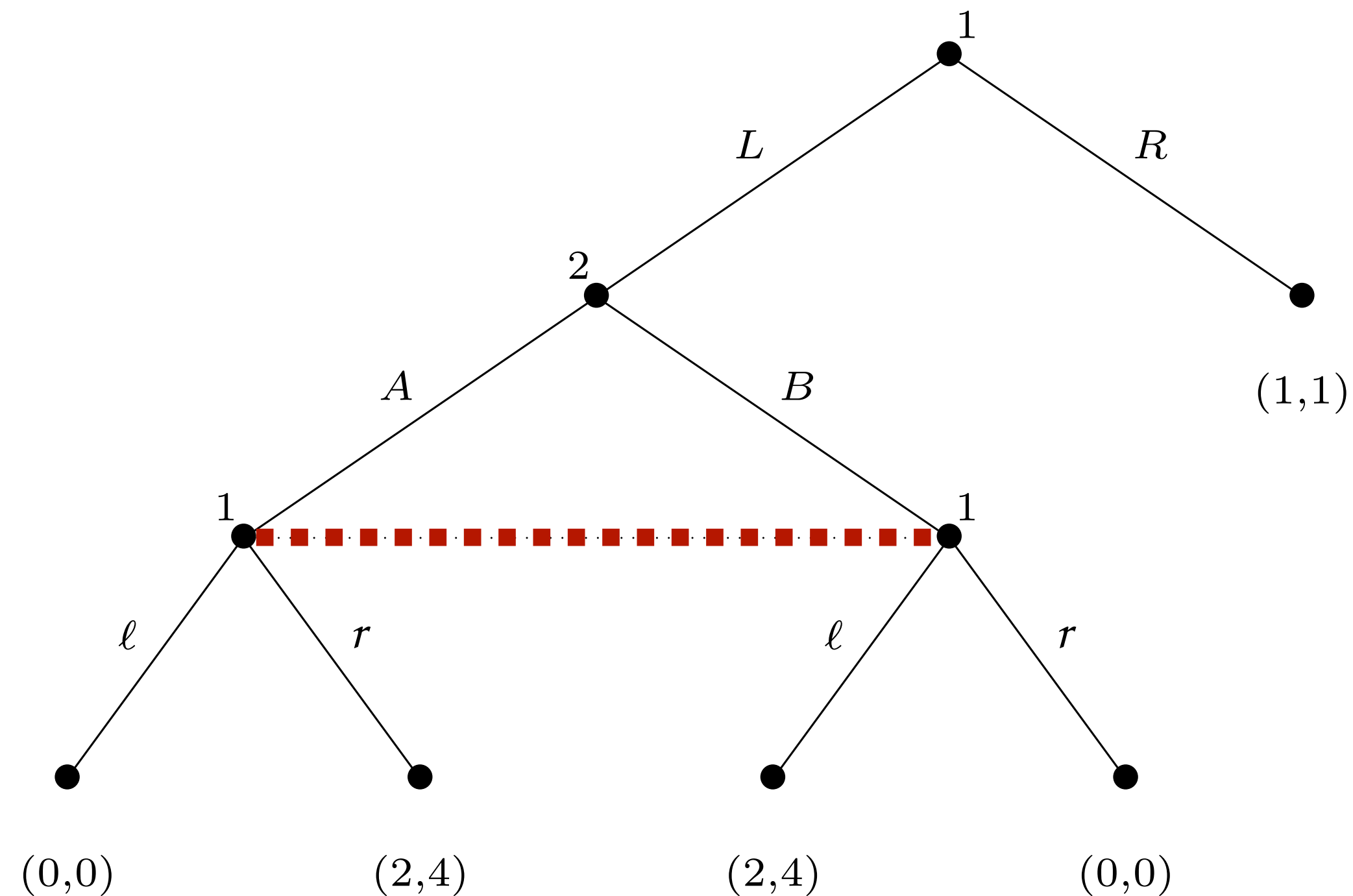
Definition:

An **imperfect information game in extensive form** is a tuple

$G = (N, A, H, Z, \chi, \rho, \sigma, u, I)$, where

- $(N, A, H, Z, \chi, \rho, \sigma, u)$ is a perfect information extensive form game, and
- $I = (I_1, \dots, I_n)$, where $I_i = (I_{i,1}, \dots, I_{i,k_i})$ is an **equivalence relation** on (i.e., partition of) $\{h \in H : \rho(h) = i\}$ with the property that $\chi(h) = \chi(h')$ and $\rho(h) = \rho(h')$ whenever there exists a j for which $h \in I_{i,j}$ and $h' \in I_{i,j}$.

Imperfect Information Extensive Form Example



- The members of the equivalence classes are sometimes called **information sets**
- Players **cannot distinguish** which **history** they are in within an information set
- **Question:** What are the information sets for each player in this game?

Pure Strategies

Question: What are the **pure strategies** in an **imperfect information** game?

Definition:

Let $G = (N, A, H, Z, \chi, \rho, \sigma, u, I)$ be an imperfect information game in extensive form. Then the **pure strategies of player i** consist of the cross product of actions available to player i at each of their **information sets**, i.e.,

$$\prod_{I_{i,j} \in I_i} \chi(h)$$

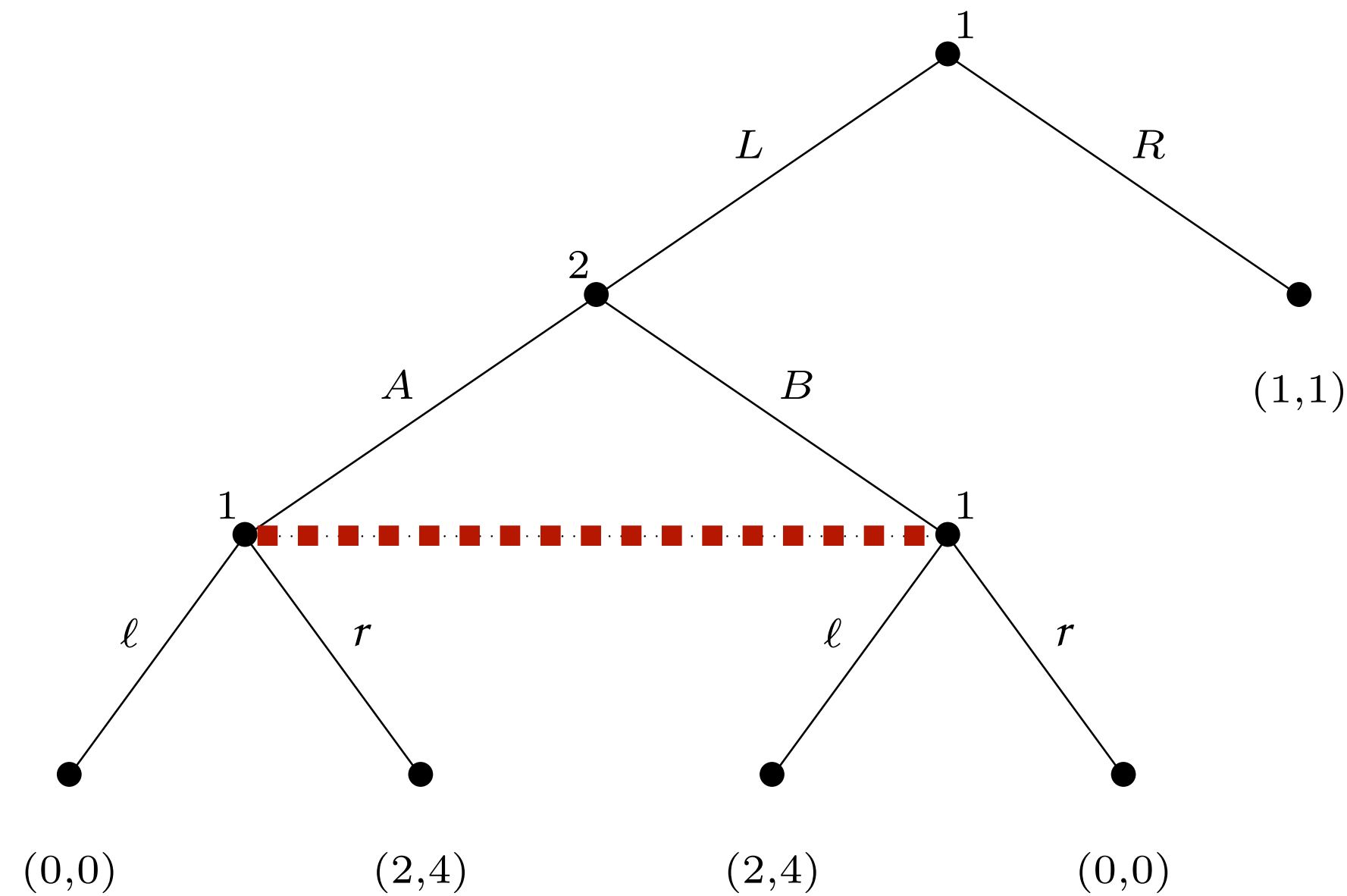
- A pure strategy associates an action with **each** information set, even those that will **never be reached**

Questions:

In an imperfect information game:

1. What are the **mixed strategies**?
2. What is a **best response**?
3. What is a **Nash equilibrium**?

Induced Normal Form



	A	B
L, ℓ	0,0	2,4
L, r	2,4	0,0
R, ℓ	1,1	1,1
R, r	1,1	1,1

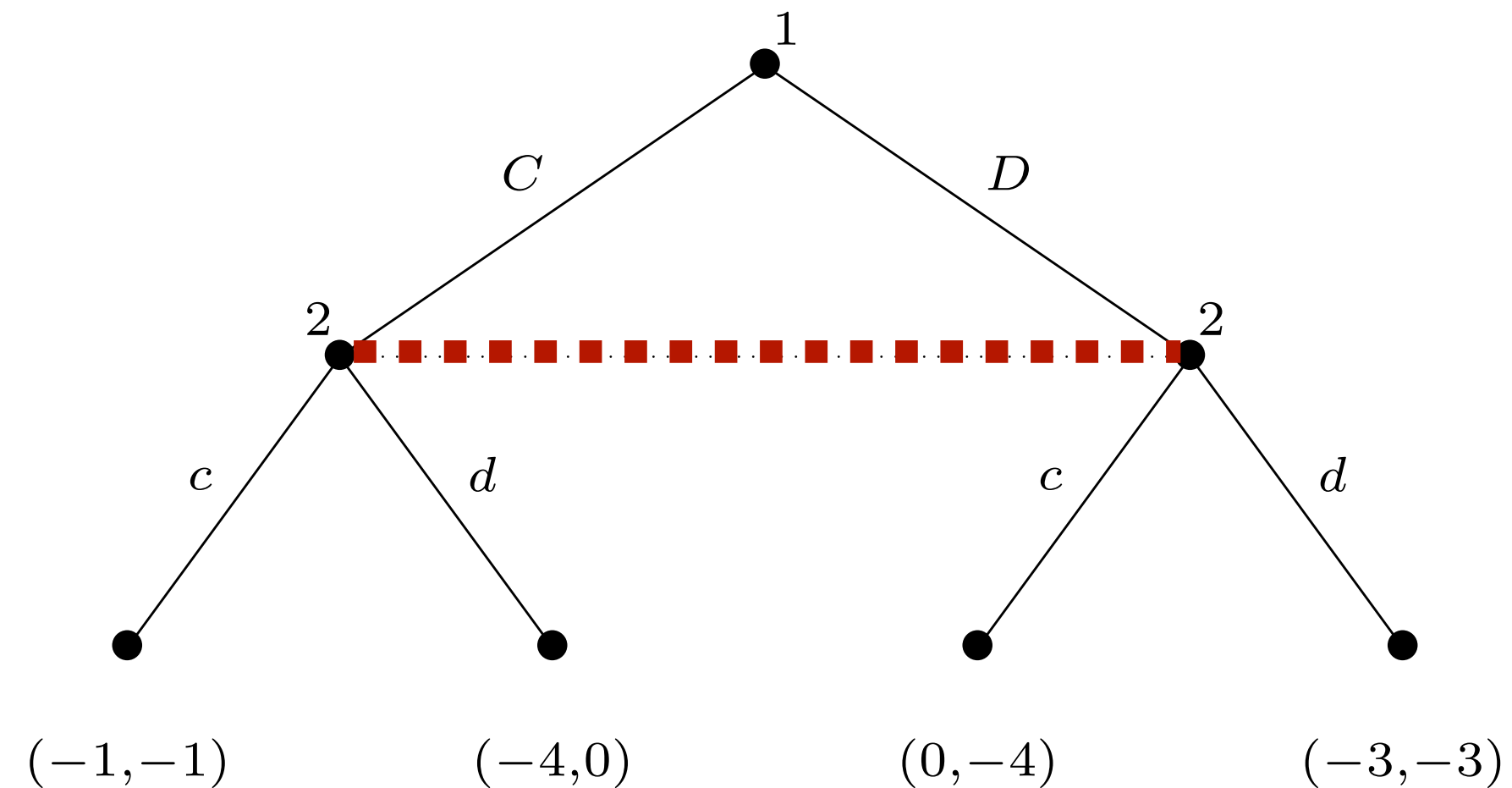
Question:

Can you represent an arbitrary **perfect information** extensive form game as an **imperfect information** game?

- Any pair of pure strategies uniquely identifies a **terminal node**, which identifies a **utility** for each agent
- We have now defined a set of **agents**, **pure strategies**, and **utility functions**
- Any extensive form game defines a corresponding **induced normal form game**

Normal to Extensive Form

	c	d
C	-1,-1	-4,0
D	0,-4	-3,-3



- Unlike perfect information games, we can go in the opposite direction and represent **any normal form game** as an **imperfect information extensive form game**
- Players can play in **any order (why?)**
- **Question:** What happens if we run this translation on the induced normal form?

Behavioural vs. Mixed Strategies

Definition:

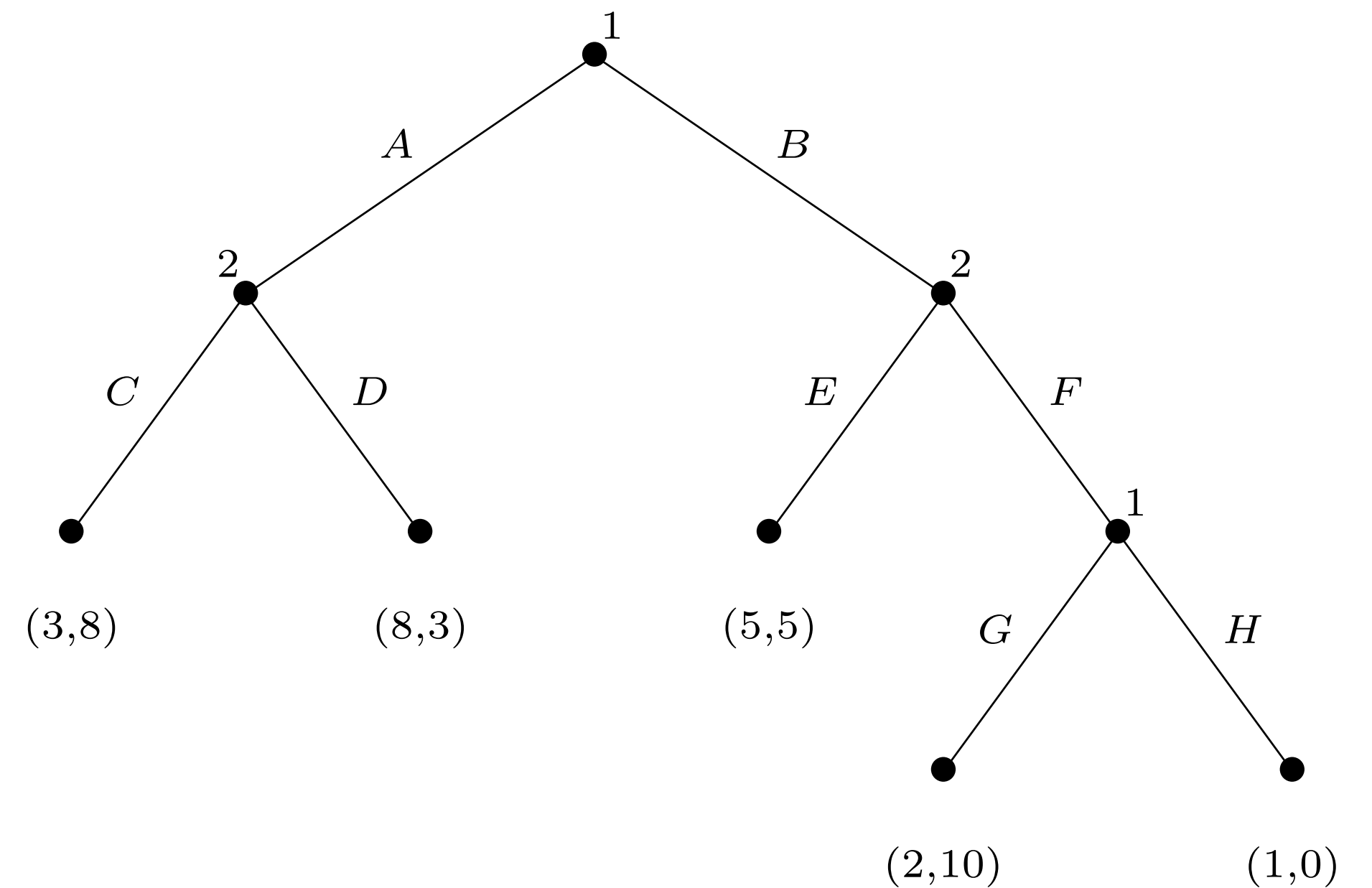
A **mixed strategy** $s_i \in \Delta(A^{I_i})$ is any distribution over an agent's **pure strategies**.

Definition:

A **behavioural strategy** $b_i \in [\Delta(A)]^{I_i}$ is a probability distribution over an agent's actions at an **information set**, which is **sampled independently** each time the agent arrives at the information set.

Behavioural vs. Mixed Example

- **Behavioural strategy:** $([.6:A, .4:B], [.6:G, .4:H])$
- **Mixed strategy:** $[.6:(A,G), .4:(B,H)]$
- **Question:** Are these strategies **equivalent**? (why?)
- **Question:** Can you construct a **mixed strategy** that is equivalent to the behavioural strategy above?
- **Question:** Can you construct a **behavioural strategy** that is equivalent to the mixed strategy above?



Perfect Recall

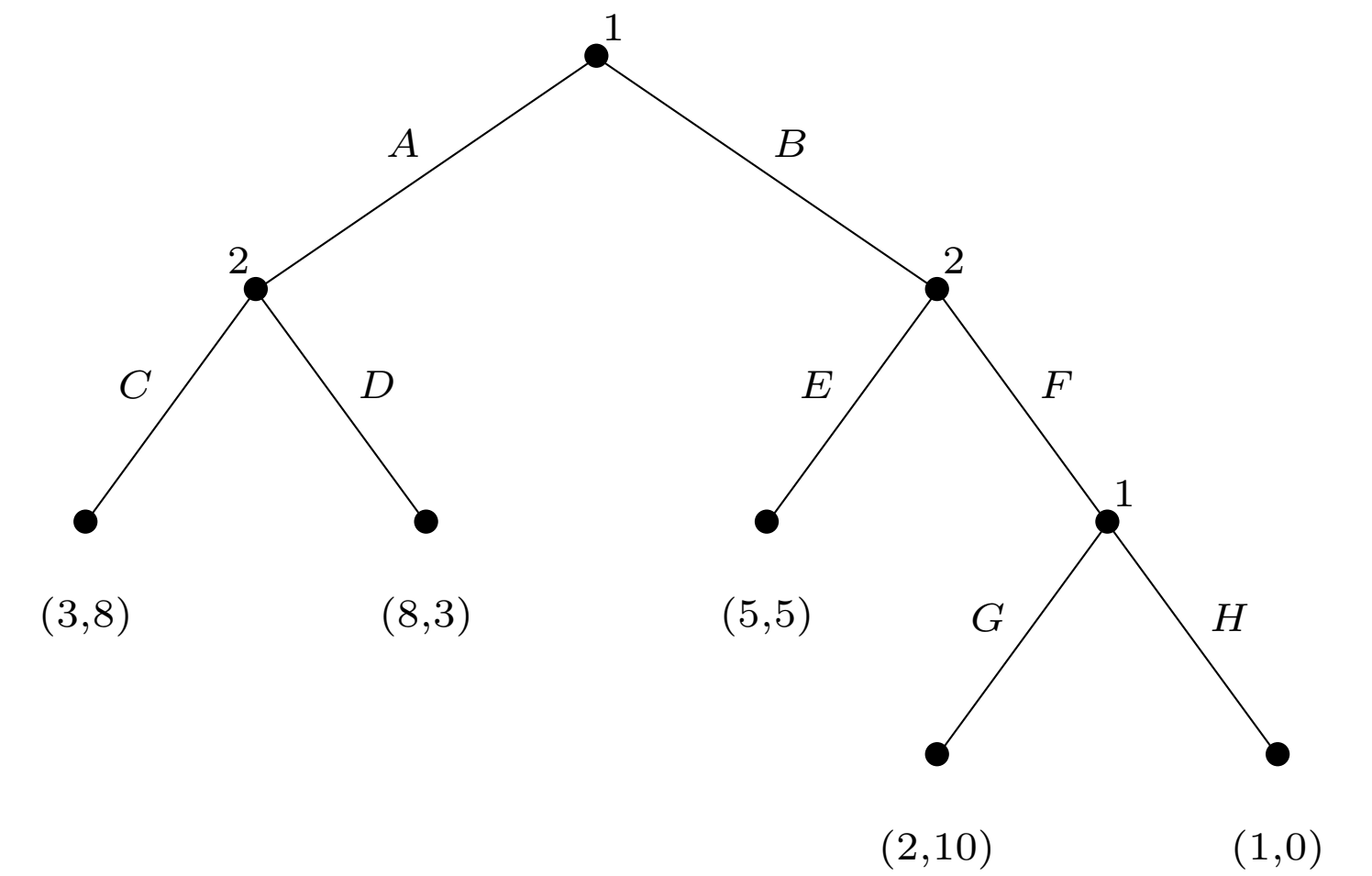
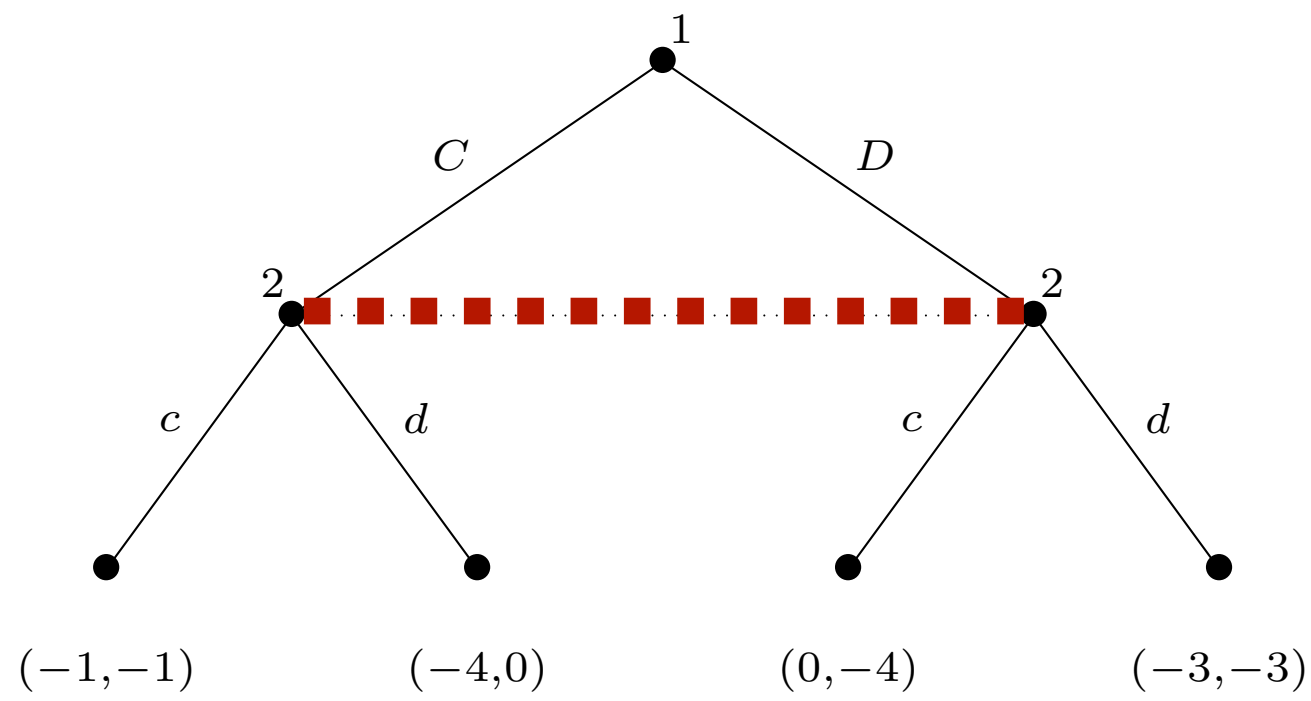
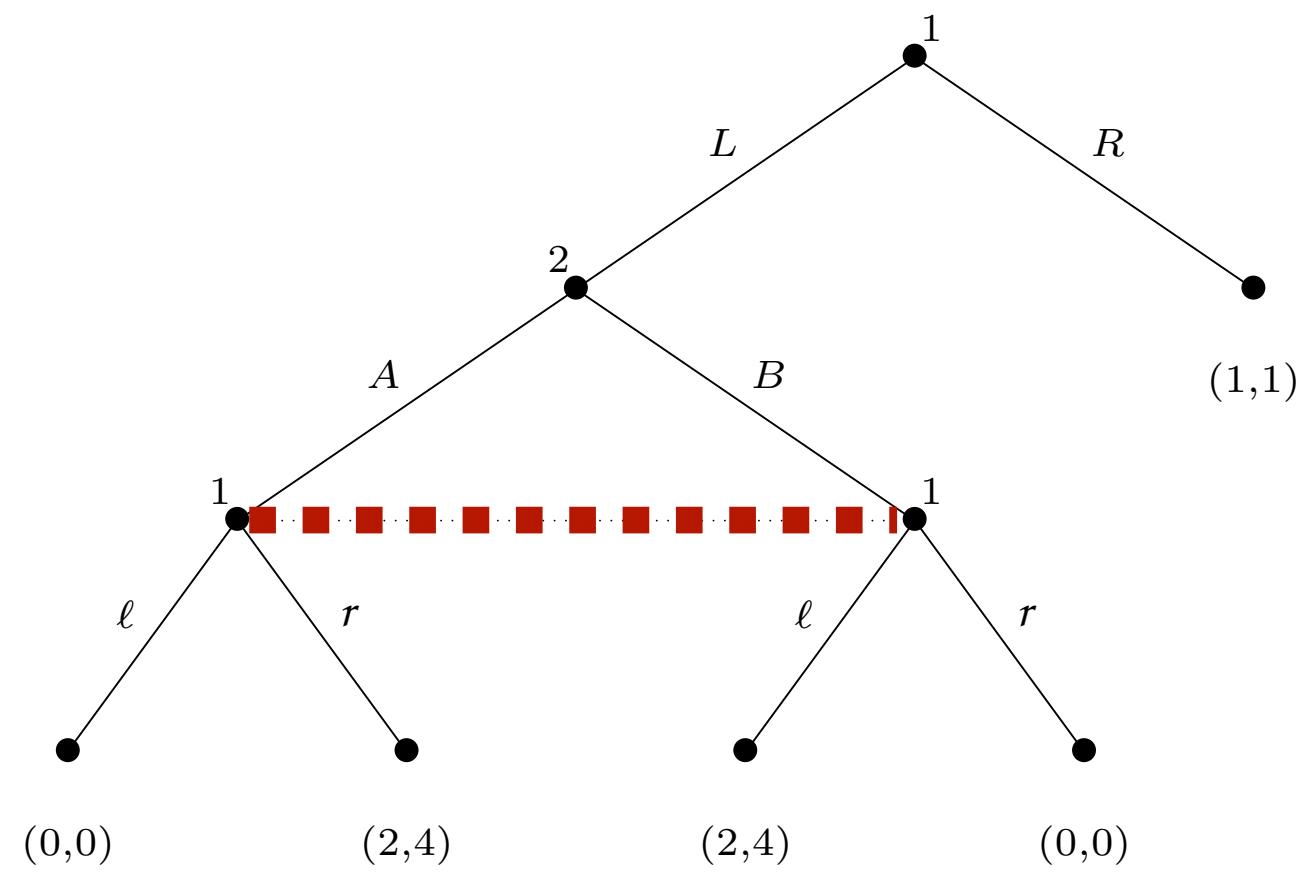
Definition:

Player i has **perfect recall** in an imperfect information game G if for any two nodes h, h' that are in the same information set for player i , for any path $h_0, a_0, h_1, a_1, \dots, h_n, h$ from the root of the game to h , and for any path $h_0, a'_0, h'_1, a'_1, \dots, h'_m, h'$ from the root of the game to h' , it must be the case that:

1. $n = m$, and
2. for all $0 \leq j \leq n$, h_j and h'_j are in the same information set, and
3. for all $0 \leq j \leq n$, if $\rho(h_j) = i$, then $a_j = a'_j$.

G is a **game of perfect recall** if every player has perfect recall in G .

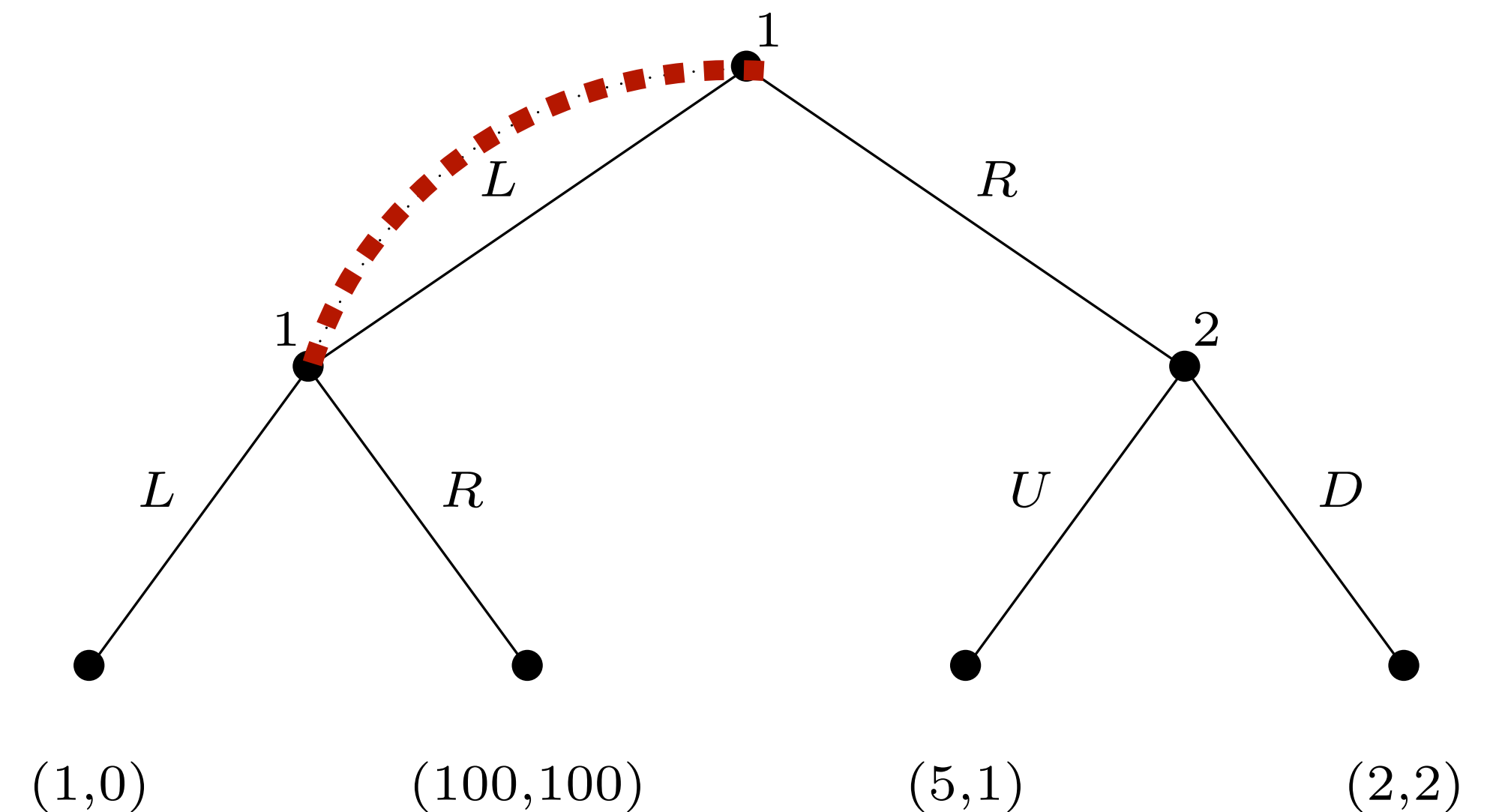
Perfect Recall Examples



Question: Which of the above games is a game of **perfect recall**?

Imperfect Recall Example

- Player 1 **doesn't remember** whether they have played L before or not. Equivalently, they visit the **same information set multiple times**
- **Question:** Can you construct a **mixed strategy** equivalent to the behavioural strategy $[\cdot 5:L, \cdot 5R]$?
- **Question:** Can you construct a **behavioural strategy** equivalent to the mixed strategy $[\cdot 5:L, \cdot 5:R]$?
- **Question:** What is the **mixed strategy equilibrium** in this game?
- **Question:** What is an **equilibrium in behavioural strategies**?



Imperfect Recall Applications

Question: When is it **useful** to model a scenario as a game of **imperfect recall**?

1. When the **actual agents** being modelled may **forget** previous history
 - Including cases where the agents strategies really are executed by **proxies**
2. As an **approximation technique**
 - E.g., **poker**: The exact cards that have been played to this point may not matter as much as some coarse grouping of which cards have been played
 - Grouping the cards into equivalence classes is a **lossy** approximation

Kuhn's Theorem

Theorem: [Kuhn, 1953]

In a game of perfect recall, any mixed strategy of a given agent can be **replaced by an equivalent behavioural strategy**, and any behavioural strategy can be **replaced by an equivalent mixed strategy**.

- Here, two strategies are **equivalent** when they induce the same probabilities on outcomes, for any fixed strategy profile (mixed or behavioural) of the other agents.

Corollary:

Restricting attention to behavioural strategies does not change the set of Nash equilibria in a game of perfect recall. (**why?**)

Computational Issues

- **Question:** Can we use **backward induction** to find an equilibrium in an imperfect information extensive form game?
- We can just use the **induced normal form** to find the equilibrium of any imperfect information game
 - But the induced normal form is **exponentially larger** than the extensive form
- Can use the **sequence form** [S&LB §5.2.3] in games of **perfect recall**:
 - **Zero-sum games:** **polynomial** in size of extensive form (i.e., exponentially faster than LP formulation on normal form)
 - **General-sum games:** **exponential** in size of extensive form (i.e., exponentially faster than converting to normal form)

Summary

- **Imperfect information extensive form games** are a model of games with sequential actions, some of which may be **hidden**
 - Histories are partitioned into **information sets**
 - Player **cannot distinguish** between histories in the same information set
- **Pure strategies** map each information set to an action
 - **Mixed strategies** are distributions over pure strategies
 - **Behavioural strategies** map each information set to a distribution over actions
 - In games of perfect recall, mixed strategies and behavioural strategies are **interchangeable**
- A player has **perfect recall** if they **never forget** anything they knew about actions so far
 - Equivalently, if they visit each information set **at most once**