## Social Choice

CMPUT 654: Modelling Human Strategic Behaviour
S\&LB §9.1-9.4

## Recap: Bayesian Games

- Epistemic types are a profile of signals that parameterize the utility functions of each agent
- Possibly correlated
- Each agent observes only their own type
- Three notions of expected utility:
- ex-ante: before observing type
- ex-interim: after observing own type
- ex-post: full type profile is known
- Solution concepts:
- Bayes-Nash equilibrium: equilibrium of induced normal form of ex-ante utilities
- Ex-post equilibrium: Agents are best-responding at every type profile (not just in expectation)


## Lecture Outline

1. Paper Bidding \& Recap
2. Aggregating Preferences
3. Voting Paradoxes
4. Arrow's Theorem

## Aggregating Preferences

- Suppose we have a collection of agents, each with individual preferences over some outcomes
- Ignore strategic reporting issues: Either The Center already knows everyone's preferences, or the agents don't lie
- How should we choose the outcome?
- More formally: Can we construct a social choice function that maps the profile of preference orderings to an outcome?
- More generally: Can we construct a social welfare function that maps the profile of preference orderings to an aggregated preference ordering?


## Inter-Agent Preference Comparisons

- Utility theory converts an ordinal preference relation into a cardinal utility function
- Can compare the strength of a single agent's preferences
- Problem: "Units" of an agent's utility function are not fixed
- Question: How can we compare the strength of two agents' preferences?
- Question: Why can't we just convert both agents' utilities to dollars (or Euros or Bitcoin or...) and compare those?


## Formal Model

Definition: A social choice function is a function $C: L^{n} \rightarrow O$, where

- $N=\{1,2, \ldots, n\}$ is a set of agents
- $O$ is a finite set of outcomes
- $L$ is the set of (non-strict) total orderings over $O$.

Definition: A social welfare function is a function $C: L^{n} \rightarrow L$, where $N, O$, and $L$ are as above.

## Notation:

We will denote $i$ 's preference order as $\succeq_{i} \in L$, and a profile of preference orders as $[\geq] \in L^{n}$.

## Non-ranking Voting Schemes

Voters need not submit a full preference ordering:

1. Plurality voting: Everyone votes for favourite outcome, choose the outcome with the most votes
2. Cumulative voting: Everyone is given $k$ votes to distribute among candidates as they like; choose the outcome with the most votes
3. Approval voting: Each agent casts a single vote for each of the outcomes that are "acceptable"; choose the outcome with the most votes.

## Ranking Voting Schemes

Every agent expresses their full preference ordering:

1. Plurality with elimination

- Everyone votes for favourite outcome
- Outcome with least votes is eliminated
- Repeat until one outcome remains

2. Borda

- Everyone assigns scores to outcome: Most-preferred gets $m-1$, next-most-preferred gets $m-2$, etc. Least-preferred outcome gets 0 .
- Outcome with highest sum of scores is chosen


## 3. Pairwise Elimination

- Define a schedule over the order in which pairs of outcomes will be compared
- For each pair, everyone chooses their favourite; least-preferred is eliminated
- Continue to next pair of non-eliminated outcomes until only one outcome remains


## Condorcet Condition

## Definition:

An outcome $o \in O$ is a Condorcet winner if $\forall o^{\prime} \in O$,

$$
\left|i \in N: o \succ_{i} o^{\prime}\right|>\left|i \in N: o^{\prime} \succ_{i} o\right| .
$$

## Definition:

A social choice function satisfies the Condorcet condition if it always selects a Condorcet winner when one exists.

- If there's one outcome that would win a pairwise vote against every other possible outcome, then perhaps we want our social choice rule to pick it
- Unfortunately, such an outcome does not always exist
- There can be cycles where A would beat B, B would beat C, C would beat A


## Paradox: Condorcet Winner

$$
\begin{aligned}
499 \text { agents: } & a>b>c \\
3 \text { agents: } & b>c>a \\
498 \text { agents: } & c>b>a
\end{aligned}
$$

- Question: Who is the Condorcet winner?
- Question: Who wins a plurality election?
- Question: Who wins under plurality with elimination?


## Paradox: <br> Sensitivity to Losing Candidate <br> 35 agents: $\quad a>c>b$ <br> 33 agents: $b>a>c$ <br> 32 agents: $\quad c>b>a$

- Question: Who wins under plurality?
- Question: Who wins under Borda?
- Question: Now drop $c$. Who wins under plurality?
- Question: After dropping $c$, who wins under Borda?


## Paradox: <br> Sensitivity to Agenda Setter <br> 35 agents: $\quad a>c>b$ <br> 33 agents: $b>a>c$ <br> 32 agents: $\quad c>b>a$

- Question: Who wins under pairwise elimination with order $a, b, c$ ?
- Question: Who wins with ordering $a, c, b$ ?
- Question: Who wins with ordering $b, c, a$ ?
- The person who sets the comparison order can cause any of the three outcomes to be picked!


## Paradox: <br> Sensitivity to Agenda Setter

$$
\begin{array}{ll}
1 \text { agent: } & b>d>c>a \\
1 \text { agent: } & a>b>d \succ c \\
1 \text { agent: } & c>a>b>d
\end{array}
$$

- Question: Who wins with ordering $a, b, c, d$ ?
- Question: What is wrong with that?


## Arrow's Theorem

These problems are not a coincidence; they affect every possible voting scheme.

## Notation:

- For this section we switch to strict total orderings $L$
- Preference ordering selected by social welfare function $W$ is $>_{W}$.


## Pareto Efficiency

## Definition:

$W$ is Pareto efficient if for any $o_{1}, o_{2} \in O$,

$$
\left(\forall i \in N: o_{1} \succ_{i} o_{2}\right) \Longrightarrow\left(o_{1} \succ_{W} o_{2}\right)
$$

- If everyone agrees that $o_{1}$ is better than $o_{2}$, then the aggregated preference order should reflect that.


## Independence of Irrelevant Alternatives

## Definition:

$W$ is independent of irrelevant alternatives if, for any $o_{1}, o_{2} \in O$ and any two preference profiles $\left[>^{\prime}\right],\left[\succ^{\prime \prime}\right] \in L$,
$\left.\left(\forall i \in N: o_{1} \succ_{i}^{\prime} o_{2} \Longleftrightarrow o_{1} \succ_{i}^{\prime \prime} o_{2}\right) \Longrightarrow\left(o_{1}\right\rangle_{W[>]} o_{2} \Longleftrightarrow o_{1} \succ_{W\left[>^{\prime \prime}\right]} o_{2}\right)$

- If every agent has the same ordering between two particular outcomes in two different preference profiles, then the social welfare function on those two profiles must order those two outcomes the same way
- The ordering between two outcomes should only depend on the agents' orderings between those outcomes, not on where any other outcomes are in the agents' orderings


## Non-Dictatorship

## Definition:

$W$ does not have a dictator if

$$
\neg i \in N: \forall[\succ] \in L^{n}: \forall o_{1}, o_{2} \in O:\left(o_{1} \succ_{i} o_{2}\right) \Longrightarrow\left(o_{1} \succ_{W} o_{2}\right)
$$

- No single agent determines the social ordering


## Arrow's Theorem

## Theorem: (Arrow, 1951)

If $|O|>2$, any social welfare function that is Pareto efficient and independent of irrelevant alternatives is dictatorial.

- Unfortunately, restricting to social choice functions instead of full social welfare functions doesn't help.


## Theorem: (Muller-Satterthwaite, 1977)

If $|O|>2$, any social choice function that is weakly Pareto efficient and monotonic is dictatorial.

## Summary

- Social choice is the study of aggregating the true preferences of a group of agents
- Social choice function: Chooses a single outcome based on preference profile
- Social welfare function: Chooses a full preference order over outcomes based on preference profile
- Well-known voting rules all lead to unfair or undesirable outcomes
- Arrow's Theorem: This is unavoidable
- Muller-Satterthwaite Theorem: ... even restricting to social choice functions

