Local Search

CMPUT 261: Introduction to Artificial Intelligence

P&M §4.7

Logistics & Assignment #1

- Assignment #1 was released last week
 See eClass
 - Due September 27 at 11:59pm
- Office hours have begun!
 - Not mandatory; for getting help from TAs
 - There are no labs for this course: You do not need to show up for your scheduled lab section
- There will be an example/practice midterm

Assignment #1 Clarification

```
def arc_cost(self, arc):
    """Returns the cost of `arc`"""
    #TODO
    return 0

def cost(self, path):
    """Returns the cost of `path`"""
    return sum( self.arc_cost(arc) for arc in path )
```

- Default implementations of arc_cost and cost assume that a path is a sequence of arcs
 - i.e., $\langle (n_1, n_2), (n_2, n_3), (n_3, n_4) \rangle$ rather than $\langle n_1, n_2, n_3, n_4 \rangle$
- This can make initializing with a zero-length path a little painful
- You are not required to use this default representation
 - Use any data structure for paths and arcs that you like

Recap

- Search problems are an extremely general encoding for choosing a sequence of actions from a start state to a goal state
- Using heuristic functions can speed this process up
 - A* search is optimal but space-intensive
 - Branch & bound depth-first search is optimal and space efficient, but needs a good starting bound
 - Iterative Deepening A* (IDA*) finds a good bound by iterative restarts (like IDS), but can be quadratically less time-efficient
- Varying the direction of search can exploit mismatches in forward and reverse branching factors

Lecture Outline

- 1. Recap & Logistics
- 2. Local Search
- 3. Hill Climbing
- 4. Randomized Algorithms

After this lecture, you should be able to:

- Implement stochastic local search and demonstrate its operation
- Implement simulated annealing and demonstrate its operation
- Identify when stochastic local search is more appropriate than graph search
- Explain the relative advantages and disadvantages of different neighbourhood specifications

Searching for Goal Nodes

Sometimes, we know how to recognize a goal node, but not how to construct one.

Example (SAT problem): Given a Boolean formula,

$$P(X) = (X_1 \lor X_2 \lor \neg X_3) \land \dots \land (\neg X_{k-2} \lor \neg X_{k-1} \lor X_k),$$

is there an assignment of truth values to the variables X_i that makes the formula true?

- State is the values of the different variables
- Easy to recognize when we've succeeded, but computing a "satisfying assignment" is NP-complete in general
- SAT is an example of a constraint satisfaction problem

Searching for Goal Nodes

We can encode SAT as a graph search problem (assignments as states, variable value changes as actions), but:

- 1. The space is too big to explore exhaustively
 - Question: How many states are there in a SAT problem with k variables?
 - Industrial SAT problems routinely have hundreds of thousands of variables
- 2. We don't care about the sequence of actions
 - Once we have a satisfying assignment, we are done
 - In fact, there isn't even a "real" set of actions; we have to make something up!

Local Search

- Idea: start from a random assignment, and then search around in the space of possible assignments
- Need not keep track of the sequence of moves that we took
- Intuitively:
 - 1. Select an assignment of a value to each variable
 - 2. Repeat:
 - (i) Select a variable to change
 - (ii) Select a new value for that variable
 - 3. until a satisfying assignment is found

Local Search Problem

Definition: Local Search Problem

- A constraint satisfaction problem: A set of variables, domains for the variables, and constraints on their joint assignment.
- Neighbours function: neighbours(n)
 - Maps from a node n to a set of "similar" nodes
- Score function: score(n)
 - Evaluates the "quality" of an assignment

Questions:

- 1. What are the nodes?
- 2. What are the goal nodes?

Neighbourhoods

- In previous graph search problems, the successor function represents states that can be reached from a given state by taking some actual action
 - In local search problems, the neighbours function is a design decision
 - We choose actions that will help us efficiently explore the space rather than trying to represent actual actions
- Usually the neighbourhood is states that differ in small ways from the current state (variable assignment)
 - E.g.: Assignments that differ in k different variables, possibly by a small amount
- Question: What might be a good neighbourhood function for SAT?

Heuristics vs. Scores

- Previously, the heuristic was optional, for improving efficiency
- In local search problems, the score function is required
 - The state space is **too big** to exhaustively explore, so uninformed search is not an option
 - Sometimes we don't even have a goal, we just want to maximize the quality of the state
- Example scores: number of satisfied clauses (in SAT); number of satisfied constraints (in CSP)
 - Note: we maximize a score (why?)

Generic Local Search Algorithm

```
Input: a constraint satisfaction problem; a neighbours function;
a score function to maximize; a stop_walk criterion
current := random assignment of values to variables
incumbent := current
repeat
  if incumbent is a satisfying assignment:
     return incumbent
  if stop_walk():
     current := new random assignment of values to variables
  else:
     select a current from neighbours(current)
  if score(current) > score(incumbent):
     incumbent := current
until termination
```

Hill Climbing

- Idea: Select the neighbour with the highest score
 - This is called an improving step
 - If no improving steps available, halt and return incumbent
- We'll move toward the best solution once we are close enough
- This algorithm is called hill climbing:
 - It seeks the highest point on the scoring function's graph
 - It moves only uphill (i.e., it makes only improving steps)

Hill Climbing Algorithm

Input: a constraint satisfaction problem; a neighbours function; a score function

```
current := random assignment of values to variables
incumbent := current
repeat
  if incumbent is a satisfying assignment:
     return incumbent
  if False:
     current := new random assignment of values to variables
  else:
     current := n \text{ from } neighbours(current) \text{ with maximum } score(n)
  if score(current) > score(incumbent):
     incumbent := current
```

Questions:

- Is hill climbing complete?
- 2. Is hill climbing optimal?

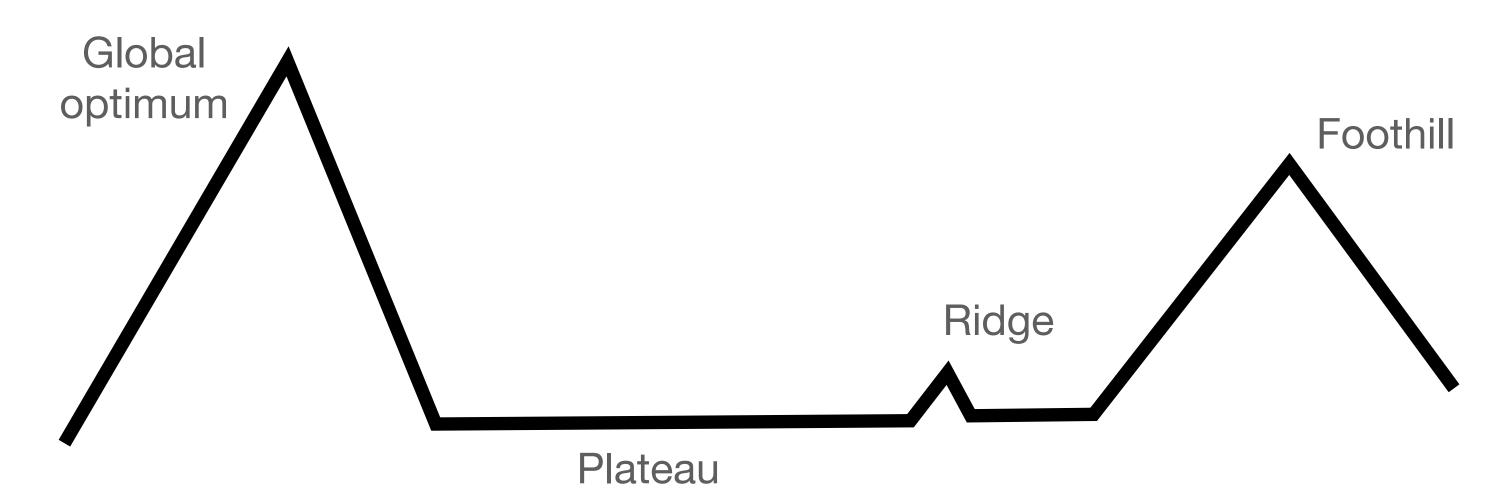
```
else:
```

return incumbent

until termination

Hill Climbing Problems

- 1. Foothills: Local maxima that are not global maxima
- 2. Plateaus: Regions of the state space where the score is uninformative
- 3. Ridges: Foothills that would not be foothills with a larger neighbourhood
- 4. **Ignorance of the global optimum:** Unless we reach a satisfying assignment, we cannot be sure that an optimum returned by local search is the **global optimum**.



Randomized Algorithms

- Adding random moves can fix some hill climbing problems
- Two main kinds of random move:
 - Random restart: Start searching from a completely random new location
 - 2. Random step: Choose a random neighbour
- Stochastic local search: Add both kinds of random moves to hill climbing

Stochastic Local Search

Input: a constraint satisfaction problem; a *neighbours* function; a *score* function to maximize; a *stop_walk* criterion; a *random_step* criterion

current := n from neighbours(current) with maximum score(n)

```
current := random assignment of values to variables
incumbent := current

repeat
  if incumbent is a satisfying assignment:
    return incumbent
  if stop_walk():
    current := new random assignment of values to variables
  else if random_step():
    current := a random element from neighbours(current)
    else:
```

if score(current) > score(incumbent):

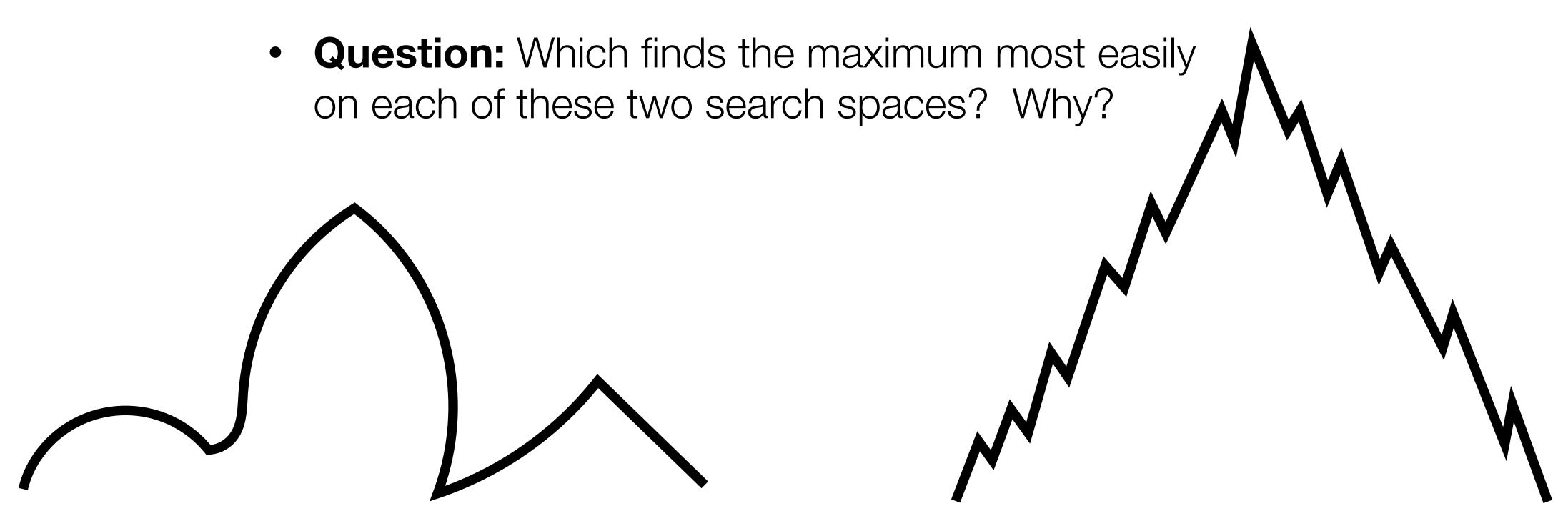
incumbent := current

Questions:

- Is stochastic local search complete?
 (Why?)
- 2. Is stochastic local search optimal? (Why?)

Two Examples

- Consider two partial algorithms:
 - 1. Hill climbing plus random restart
 - 2. Hill climbing plus random steps



Simulated Annealing

- Idea: Start out by searching pretty randomly, but become more directed
 - Intuition: Move to a good neighbourhood quickly, then search intensively in that neighbourhood
- Maintain a "temperature" T
- Choose new nodes more randomly at higher temperatures;
 Gradually decrease the temperature (according to a cooling schedule)
- At each step:
 - 1. Randomly choose a neighbour *new*
 - 2. If score(new) > score(current), always accept (i.e., assign to current)
 - 3. Else, accept with probability

$$\rho[(score(new)-score(current))/T]$$

Simulated Annealing cont.

$$e^{[(score(new)-score(current))/T]}$$

- Worse score(new) means lower acceptance probability
- Always negative (why?)

- Higher T makes negative value smaller
- Higher acceptance probability
- Small neighbourhoods are good, because they are more efficient to search
- Large neighbourhoods are good, because they are more likely to contain an improvement
- Simulated annealing allows for a large neighbourhood and efficient searching
 - You don't have to generate the whole neighbourhood, just randomly construct a single neighbour

Summary

- For some problems, we only care about finding a goal node, not the actions
 we took to find it
- Local search: Look for goal states by iteratively moving from a current state to a neighbouring state
 - Hill climbing: Always move to the highest-score neighbour
 - Random step: Sometimes choose a random neighbour
 - Random restart: Sometimes start again from an entirely random state
 - Simulated annealing: Random moves start very random, become more greedy over time