## Social Choice

CMPUT 654: Modelling Human Strategic Behaviour

S&LB §9.1-9.4

# Logistics

- Assignment #1: Late deadline is tonight at 11:59pm
- Paper bidding:
  - Next lecture we will assign papers to presenters
  - Please look over the abstracts to figure out your own preferences!
  - And/or bring a different paper that you'd really like to present
  - If you want to present a different paper:
    - You will still need to bid for a timeslot
    - You will need to send it before class for approval

#### Part 2: Behavioural game theory

wk	Date	Topic	Readings & Milestones
1	Tue, Feb 27	Behavioral economics intro	Kahneman & Tversky (1979) Conlisk (1989) Project talk scheduling
1	Thu, Feb 29	Single-shot interactions	Camerer, Ho, and Chong (2004) McKelvey & Palfrey (1995) Wright & Leyton-Brown (2017)
2	Tue, Mar 5	Salience and focal points	Crawford & Iriberri (2007) Burchardi and Penczynski (2014) Wright & Leyton-Brown (2019)
2	Thu, Mar 7	Fairness and social preferences	Kahneman, Knetsch, and Thaler (1986) Gal et al. (2017) Assignment 2 due
3	Tue, Mar 12	Reasoning about sequential interactions	Li (2017) Nagel & Saitto (2023) Zinkevich et al. (2007) Survey outline due
3	Thu, Mar 14	No-regret learning	Hart & Mas-Colell (2000) Nekipelov, Syrgkanis, and Tardos (2015) Morrill et al. (2021)
4	Tue, Mar 19	Stackelberg equilibrium	Kiekintveld et al. (2009) Deng, Schneider, and Sivan (2019) Brown et al. (2023)
4	Thu, Mar 21	Behavioral finance	Benartzi & Thaler (1995) Khaw et al. (2017) Rabin (2000)

## Recap: Bayesian Games

- Epistemic types are a profile of signals that parameterize the utility functions of each agent
  - Possibly correlated
  - Each agent observes only their own type
- Three notions of expected utility:
  - ex-ante: before observing type
  - ex-interim: after observing own type
  - ex-post: full type profile is known
- Solution concepts:
  - Bayes-Nash equilibrium: equilibrium of induced normal form of ex-ante utilities
  - **Ex-post equilibrium:** Agents are best-responding at **every type profile** (not just in expectation)

### Lecture Outline

- 1. Logistics & Recap
- 2. Aggregating Preferences
- 3. Voting Paradoxes
- 4. Arrow's Theorem

# Familiar-Looking Question

- Consider the lottery [.9999: \$1, .0001: \$10,000,000] Expected value: \$1000.9999
- Question: Can a rational expected utility maximizer with monotonic value for money turn down a ticket to this lottery that costs \$100?

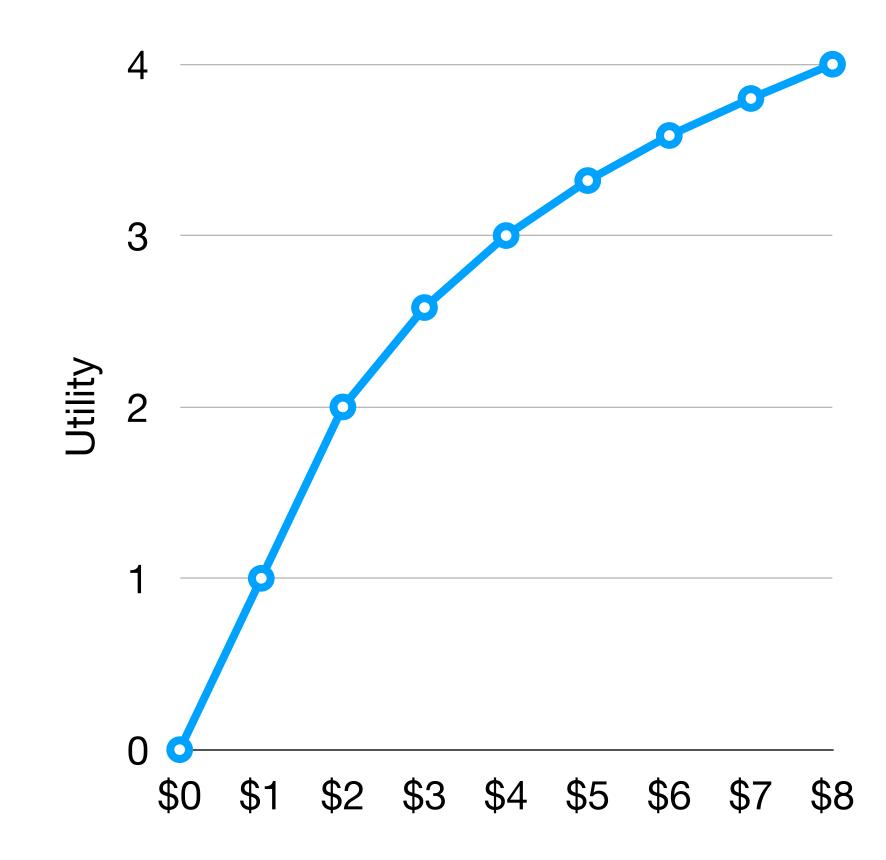
### Risk Attitudes

- Rational agents are required to maximize their expected utility
  - This is distinct from the expected value of some outcome (e.g., money)
- Rational agents can have a utility for money that is:
  - Risk-averse: Value of a lottery is lower than expected value
  - Risk-neutral: Value of a lottery is exactly expected value
  - Risk-seeking: Value of a lottery is higher than expected value

# Risk Aversion Example

- Example:  $\ell = [.5:\$0,.5:\$8]$  has expected value of \$4
- Question: What is the expected utility of ℓ?
- Expectations are taken over **outcomes**:  $\mathbb{E}[u(o)] \neq u(\mathbb{E}[o])$

$$u(\ell) = \mathbb{E}[u(o)] = .5u(\$0) + .5u(\$8)$$
$$= 2 = .5(0) + .5(4)$$
$$\neq 3 = u(\$4) = u(\mathbb{E}[o])$$



# Risk Attitudes and Utility for Money

- Concave utility for money 

   risk aversion
- Linear utility for money 

   risk neutral
- Convex utility for money 

   risk seeking

## Aggregating Preferences

- Suppose we have a collection of agents, each with individual preferences over some outcomes
  - Ignore strategic reporting issues: Either **The Center** already knows everyone's preferences, or the agents don't lie
- How should we choose the outcome?
- More formally: Can we construct a social choice function that maps the profile of preference orderings to an outcome?
- More generally: Can we construct a social welfare function that maps the profile of preference orderings to an aggregated preference ordering?

# Inter-Agent Preference Comparisons

- Utility theory converts an ordinal preference relation into a cardinal utility function
  - Can compare the strength of a single agent's preferences
- Problem: "Units" of an agent's utility function are not fixed
- Question: How can we compare the strength of two agents' preferences?
- **Question:** Why can't we just convert both agents' utilities to **dollars** (or Euros or Bitcoin or...) and compare *those*?

### Formal Model

**Definition:** A social choice function is a function  $C:L^n\to O$ , where

- $N = \{1, 2, ..., n\}$  is a set of **agents**
- O is a finite set of outcomes
- L is the set of (non-strict) total orderings over O.

**Definition:** A social welfare function is a function  $C:L^n\to L$ , where N,O, and L are as above.

#### **Notation:**

We will denote i's preference order as  $\geq_i \in L$ , and a profile of preference orders as  $[\geq] \in L^n$ .

## Non-ranking Voting Schemes

Voters need not submit a full preference ordering:

- 1. **Plurality voting:** Everyone votes for favourite outcome, choose the outcome with the most votes
- 2. Cumulative voting: Everyone is given k votes to distribute among candidates as they like; choose the outcome with the most votes
- 3. **Approval voting:** Each agent casts a single vote for each of the outcomes that are "acceptable"; choose the outcome with the most votes.

# Ranking Voting Schemes

Every agent expresses their full preference ordering:

#### 1. Plurality with elimination

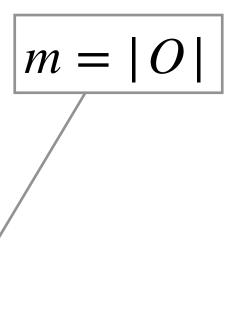
- Everyone votes for favourite outcome
- Outcome with least votes is eliminated
- Repeat until one outcome remains

#### 2. **Borda**

- Everyone assigns scores to outcome: Most-preferred gets m-1, next-most-preferred gets m-2, etc. Least-preferred outcome gets 0.
- Outcome with highest sum of scores is chosen

#### 3. Pairwise Elimination

- Define a schedule over the order in which pairs of outcomes will be compared
- For each pair, everyone chooses their favourite; least-preferred is eliminated
- Continue to next pair of non-eliminated outcomes until only one outcome remains



### Condorcet Condition

#### **Definition:**

An outcome  $o \in O$  is a Condorcet winner if  $\forall o' \in O$ ,

$$|i \in N: o >_i o'| > |i \in N: o' >_i o|$$
.

#### **Definition:**

A social choice function satisfies the **Condorcet condition** if it always selects a Condorcet winner when one exists.

- If there's one outcome that would win a pairwise vote against every other possible outcome, then perhaps we want our social choice rule to pick it
- Unfortunately, such an outcome does not always exist
  - There can be cycles where A would beat B, B would beat C, C would beat A

# Paradox: Condorcet Winner

499 agents: a > b > c

3 agents: b > c > a

498 agents: c > b > a

- Question: Who is the Condorcet winner?
- Question: Who wins a plurality election?
- Question: Who wins under plurality with elimination?

# Paradox: Sensitivity to Losing Candidate

- 35 agents: a > c > b
- 33 agents: b > a > c
- 32 agents: c > b > a
- Question: Who wins under plurality?
- Question: Who wins under Borda?
- Question: Now drop c. Who wins under plurality?
- Question: After dropping c, who wins under Borda?

# Paradox: Sensitivity to Agenda Setter

35 agents: a > c > b

33 agents: b > a > c

32 agents: c > b > a

- Question: Who wins under pairwise elimination with order a,b,c?
- Question: Who wins with ordering a,c,b?
- Question: Who wins with ordering b,c,a?
- The person who sets the comparison order can cause **any** of the three outcomes to be picked!

# Paradox: Sensitivity to Agenda Setter

- 1 agent: b > d > c > a
- 1 agent: a > b > d > c
- 1 agent: c > a > b > d

- Question: Who wins with ordering a,b,c,d?
- Question: What is wrong with that?

### Arrow's Theorem

These problems are not a coincidence; they affect every possible voting scheme.

#### **Notation:**

- ullet For this section we switch to strict total orderings L
- Preference ordering selected by social welfare function W is  $\succ_W$ .

## Pareto Efficiency

#### **Definition:**

W is Pareto efficient if for any  $o_1, o_2 \in O$ ,

$$(\forall i \in N : o_1 \succ_i o_2) \implies (o_1 \succ_W o_2).$$

• If everyone agrees that  $o_1$  is better than  $o_2$ , then the aggregated preference order should reflect that.

## Independence of Irrelevant Alternatives

#### **Definition:**

W is independent of irrelevant alternatives if, for any  $o_1, o_2 \in O$  and any two preference profiles  $[\succ'], [\succ''] \in L$ ,

$$(\forall i \in N : o_1 \succ_i' o_2 \iff o_1 \succ_i'' o_2) \implies (o_1 \succ_{W[\succ']} o_2 \iff o_1 \succ_{W[\succ'']} o_2)$$

- If every agent has the **same ordering** between **two particular outcomes** in two different preference profiles, then the social welfare function on those two profiles must order those two outcomes the same way
- The ordering between two outcomes should only depend on the agents' orderings between those outcomes, not on where any other outcomes are in the agents' orderings

## Non-Dictatorship

#### **Definition:**

W does not have a dictator if

$$\neg i \in N : \forall [>] \in L^n : \forall o_1, o_2 \in O : (o_1 >_i o_2) \implies (o_1 >_W o_2)$$

No single agent determines the social ordering

### Arrow's Theorem

Theorem: (Arrow, 1951)

If |O| > 2, any social welfare function that is Pareto efficient and independent of irrelevant alternatives is dictatorial.

• Unfortunately, restricting to social choice functions instead of full social welfare functions doesn't help.

Theorem: (Muller-Satterthwaite, 1977)

If |O| > 2, any social choice function that is weakly Pareto efficient and monotonic is dictatorial.

# Summary

- Social choice is the study of aggregating the true preferences of a group of agents
  - Social choice function: Chooses a single outcome based on preference profile
  - Social welfare function: Chooses a full preference order over outcomes based on preference profile
- Well-known voting rules all lead to unfair or undesirable outcomes
  - Arrow's Theorem: This is unavoidable
  - Muller-Satterthwaite Theorem: ... even restricting to social choice functions