Game Theory for Sequential Interactions

CMPUT 366: Intelligent Systems

S&LB §5.0-5.2.2

Lecture Outline

- 1. Recap
- 2. Perfect Information Games
- 3. Backward Induction
- 4. Imperfect Information Games

Recap: Game Theory

- Game theory studies the interactions of rational agents
 - Canonical representation is the normal form game
- Game theory uses solution concepts rather than optimal behaviour
 - "Optimal behaviour" is not clear-cut in multiagent settings Soccer
 - Pareto optimal: no agent can be made better off without making some other agent worse off
 - Nash equilibrium: no agent regrets their strategy given the choice of the other agents' strategies

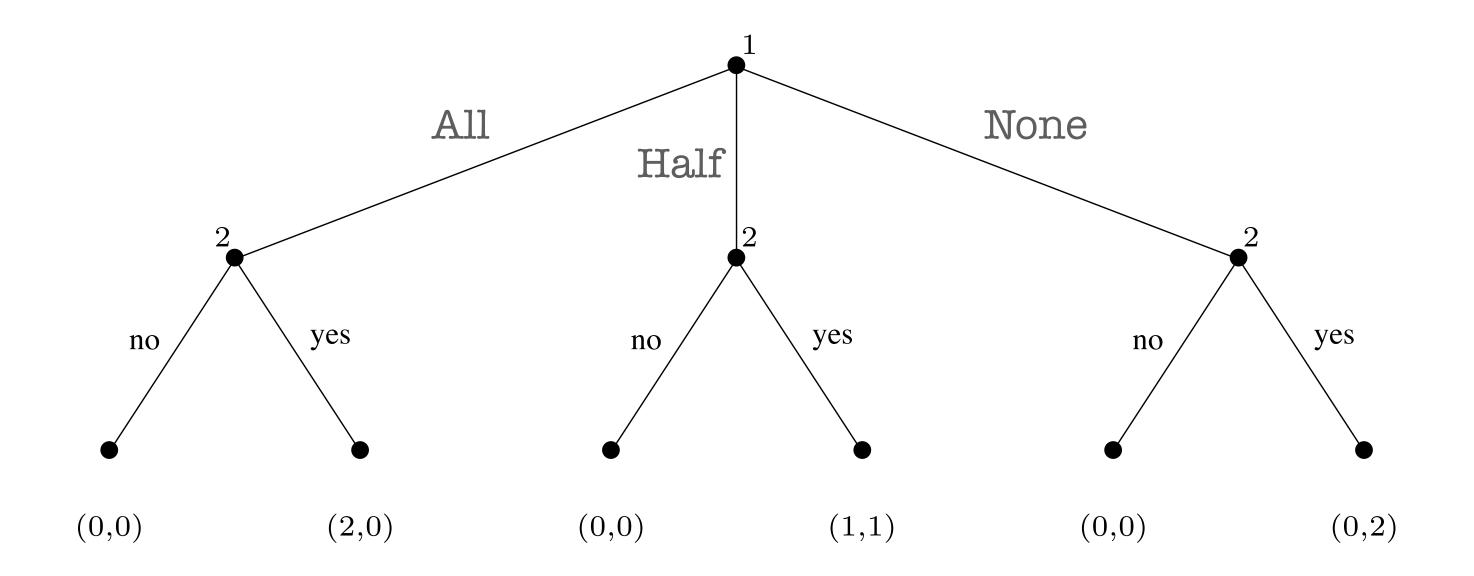
Ballet Soccer

2, 1	0, 0
0, 0	1, 2

Ballet

Extensive Form Games

- Normal form games don't have any notion of sequence: all actions happen simultaneously
- The extensive form is a game representation that explicitly includes temporal structure (i.e., a game tree)



Perfect Information

There are two kinds of extensive form game:

- 1. **Perfect information:** Every agent **sees all actions** of the other players (including "**Nature**")
 - e.g.: Chess, checkers, Pandemic
- 2. Imperfect information: Some actions are hidden
 - Players may not know exactly where they are in the tree
 - e.g.: Poker, rummy, Scrabble

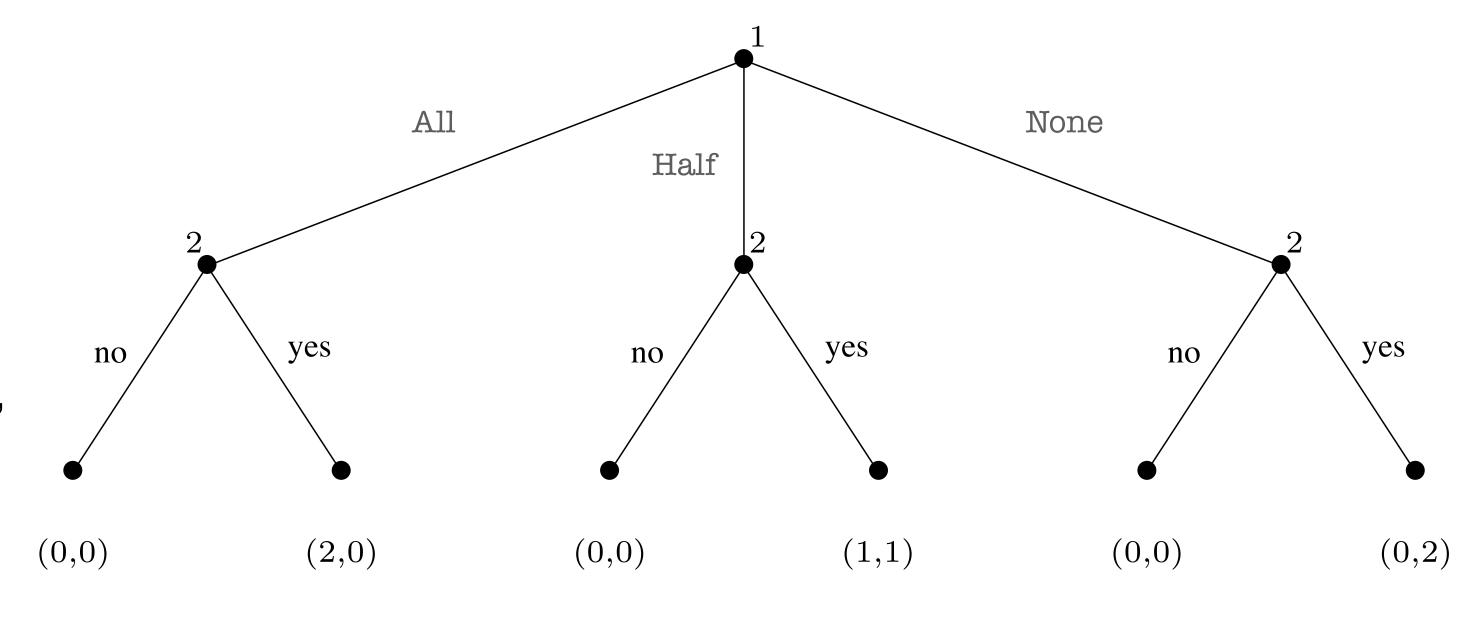
Perfect Information Extensive Form Game

Definition:

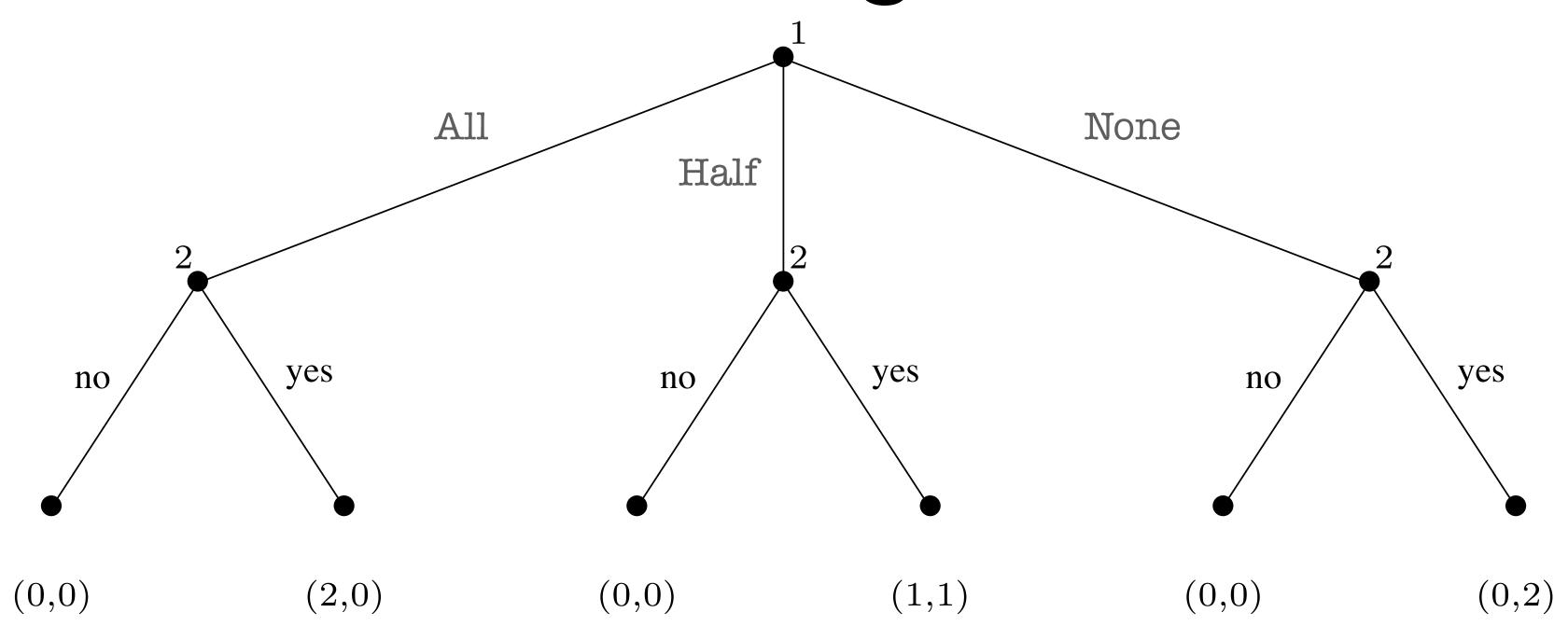
A finite perfect-information game in extensive form is a tuple

$$G = (N, A, H, Z, \chi, \rho, \sigma, u)$$
, where

- N is a set of n players,
- A is a single set of actions,
- H is a set of nonterminal choice nodes,
- Z is a set of **terminal nodes** (disjoint from H),
- $\chi: H \to 2^A$ is the action function,
- $\rho: H \to N$ is the player function,
- $\sigma: H \times A \rightarrow H \cup Z$ is the successor function,
- $u = (u_1, u_2, ..., u_n)$ is a utility function for each player, $u_i : Z \to \mathbb{R}$



Fun Game: The Sharing Game



- Two siblings must decide how to share two \$100 coins
- Sibling 1 suggests a division, then sibling 2 accepts or rejects
 - If rejected, nobody gets any coins.

Pure Strategies

Question: What are the pure strategies in an extensive form game?

Definition:

Let $G = (N, A, H, Z, \chi, \rho, \sigma, u)$ be a perfect information game in extensive form. Then the **pure strategies of player** i consist of the cross product of actions available to player i at each of their choice nodes, i.e.,

$$\prod_{h \in H \mid \rho(h) = i} \chi(h)$$

 A pure strategy associates an action with each choice node, even those that will never be reached

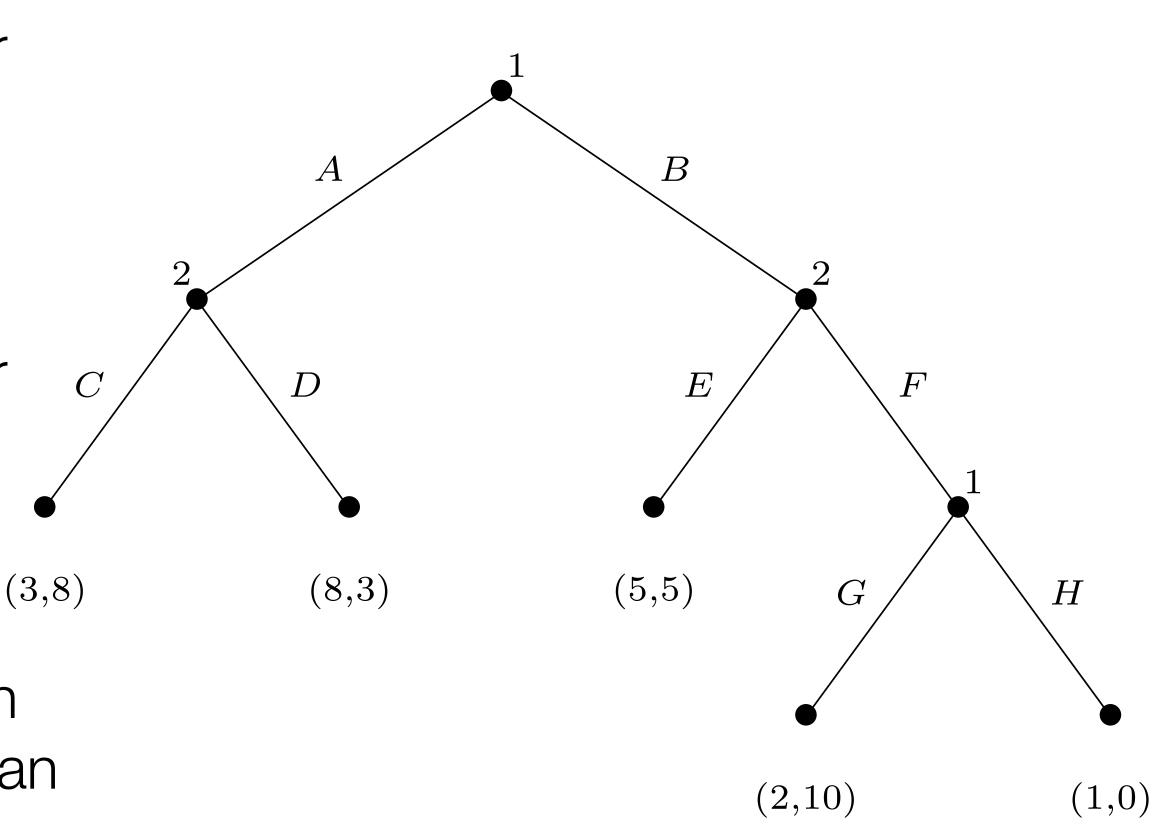
Pure Strategies Example

Question: What are the pure strategies for player 2?

• $\{(C, E), (C, F), (D, E), (D, F)\}$

Question: What are the pure strategies for player 1?

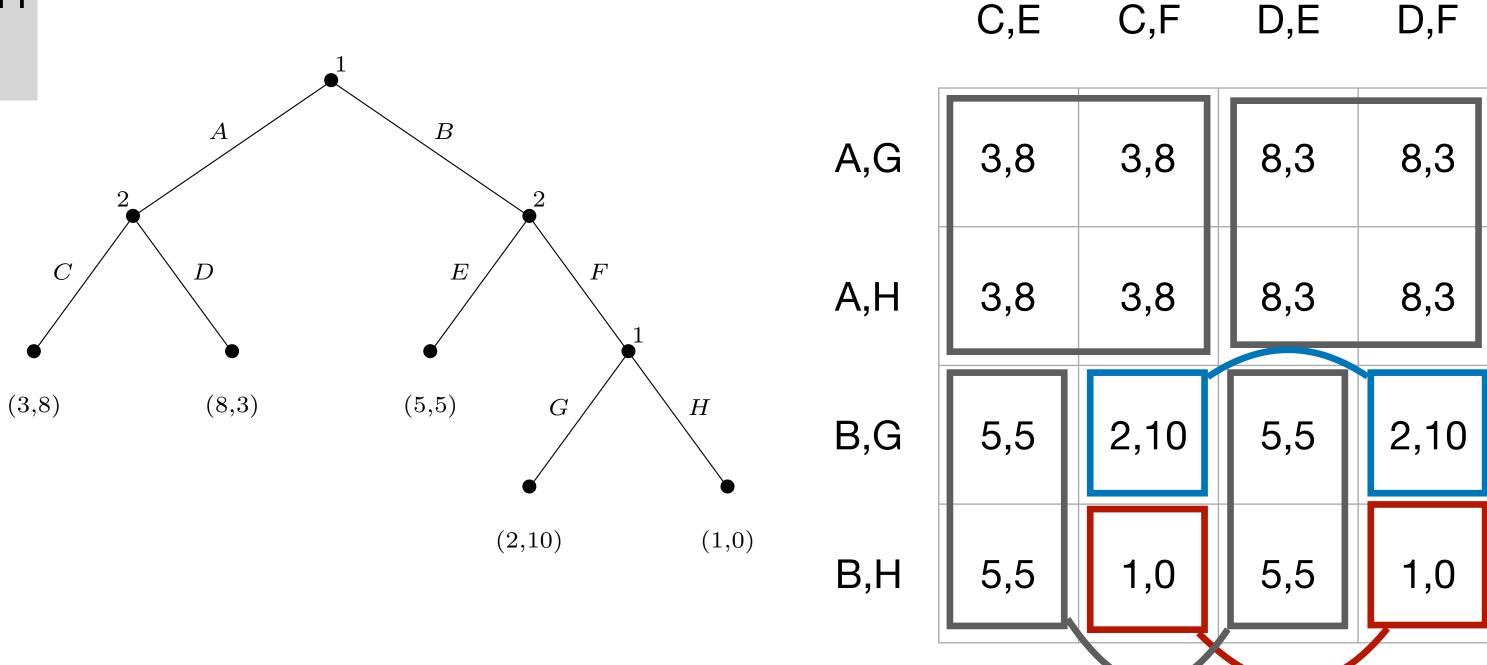
- $\{(A,G),(A,H),(B,G),(B,H)\}$
- Note that these associate an action with the second choice node even when it can never be reached



Induced Normal Form

Question:

Which representation is more **compact**?



- Any pair of pure strategies uniquely identifies a terminal node, which identifies a utility for each agent (why?)
- We have now defined a set of agents, pure strategies, and utility functions
- Any extensive form game defines a corresponding induced normal form game

Reusing Old Definitions

- We can plug our new definition of pure strategy into our existing definitions for:
 - Mixed strategy
 - Best response
 - Nash equilibrium (both pure and mixed strategy)

Question:

What is the definition of a mixed strategy in an extensive form game?

Pure Strategy Nash Equilibria

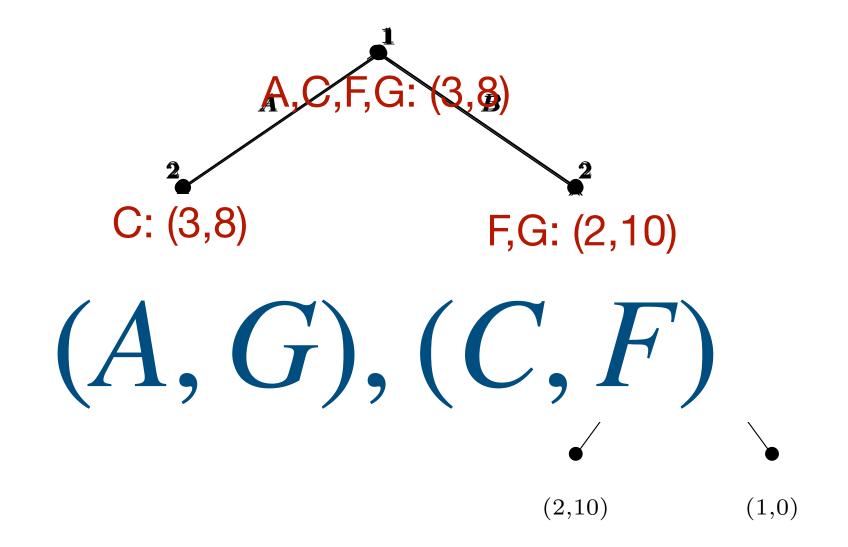
Theorem: [Zermelo, 1913]

Every finite perfect-information game in extensive form has at least one pure strategy Nash equilibrium.

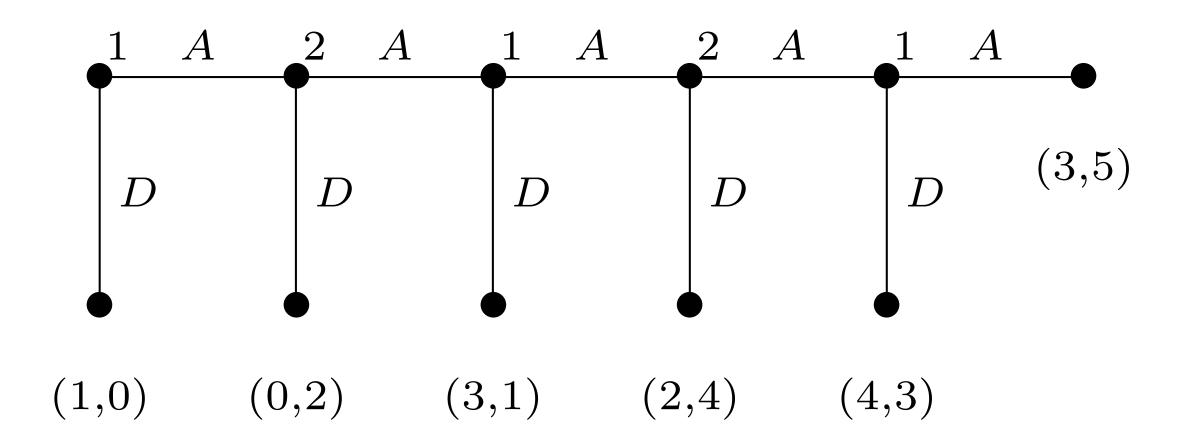
- Starting from the bottom of the tree, no agent needs to **randomize**, because they already know the best response
- There might be **multiple** pure strategy Nash equilibria in cases where an agent has multiple best responses at a **single choice node**

Backward Induction

- Backward induction is a straightforward algorithm that is guaranteed to compute a pure strategy Nash equilibrium.
- Idea: Replace subgames lower in the tree with their equilibrium values

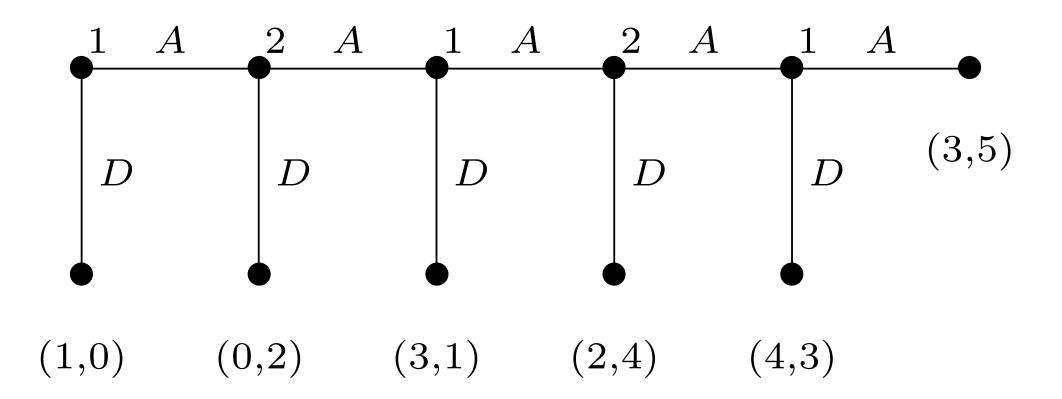


Fun Game: Centipede



- At each stage, one of the players can go Across or Down
- If they go Down, the game ends.

Backward Induction Criticism



- The unique equilibrium is for each player to play Down at the first opportunity.
- Empirically, this is not how real people tend to play!
- Theoretically, what should you do if you arrive at an off-path node?
 - How do you update your beliefs to account for this probability 0 event?
 - If player 1 knows that you will update your beliefs in a way that causes you not to play **Down**, then playing **Down** is no longer their only rational choice...

Imperfect Information, informally

- Perfect information games model sequential actions that are observed by all players
 - Randomness can be modelled by a special Nature player with constant utility and known mixed strategy
- But many games involve hidden actions
 - Cribbage, poker, Scrabble
 - Sometimes actions of the **players** are hidden, sometimes **Nature's** actions are hidden, sometimes both
- Imperfect information extensive form games are a model of games with sequential actions, some of which may be hidden

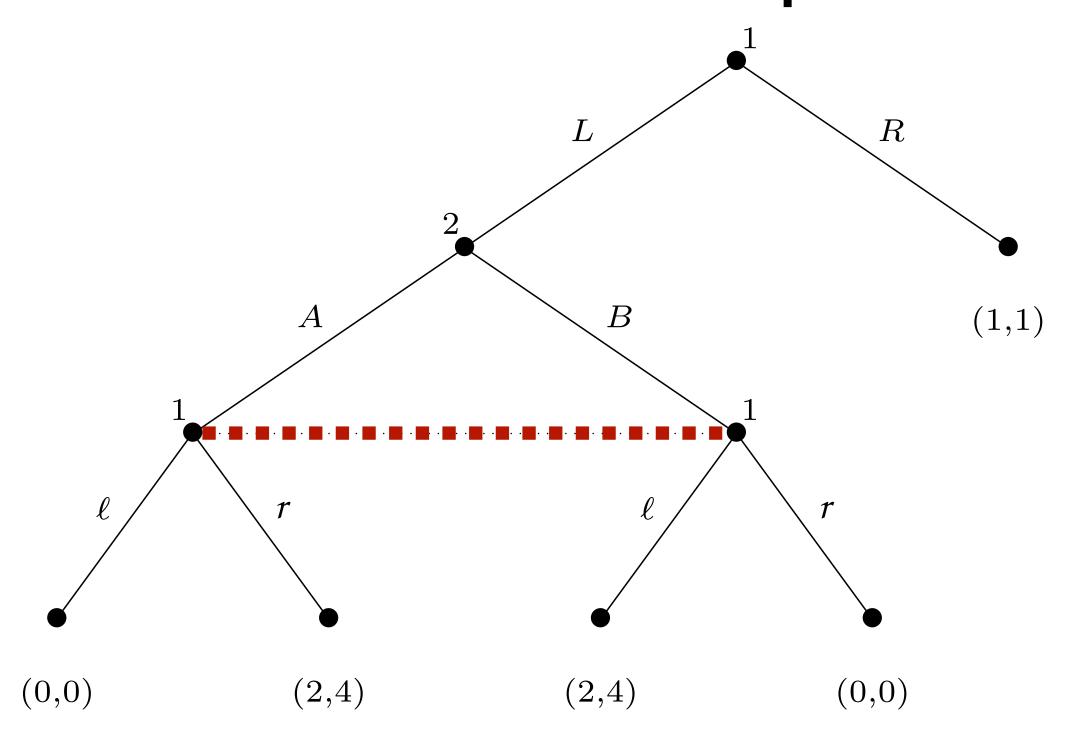
Imperfect Information Extensive Form Game

Definition:

An imperfect information game in extensive form is a tuple $G = (N, A, H, Z, \chi, \rho, \sigma, u, I)$, where

- $(N, A, H, Z, \chi, \rho, \sigma, u)$ is a perfect information extensive form game, and
- $I=(I_1,\ldots,I_n)$, where $I_i=(I_{i,1},\ldots,I_{i,k_i})$ is an **equivalence relation** on (i.e., partition of) $\{h\in H: \rho(h)=i\}$ with the property that $\chi(h)=\chi(h')$ and $\rho(h)=\rho(h')$ whenever there exists a $j\in N$ for which $h\in I_{i,j}$ and $h'\in I_{i,j}$.

Imperfect Information Extensive Form Example



- The members of the equivalence classes are also called information sets
- Players cannot distinguish which history they are in within an information set
- Question: What are the information sets for each player in this game?

Pure Strategies

Question: What are the pure strategies in an imperfect information extensive-form game?

Definition:

Let $G = (N, A, H, Z, \chi, \rho, \sigma, u, I)$ be an imperfect information game in extensive form. Then the **pure strategies of player** i consist of the cross product of actions available to player i at each of their **information sets**, i.e.,

$$\prod_{I_{i,j}\in I_i} \chi(h)$$

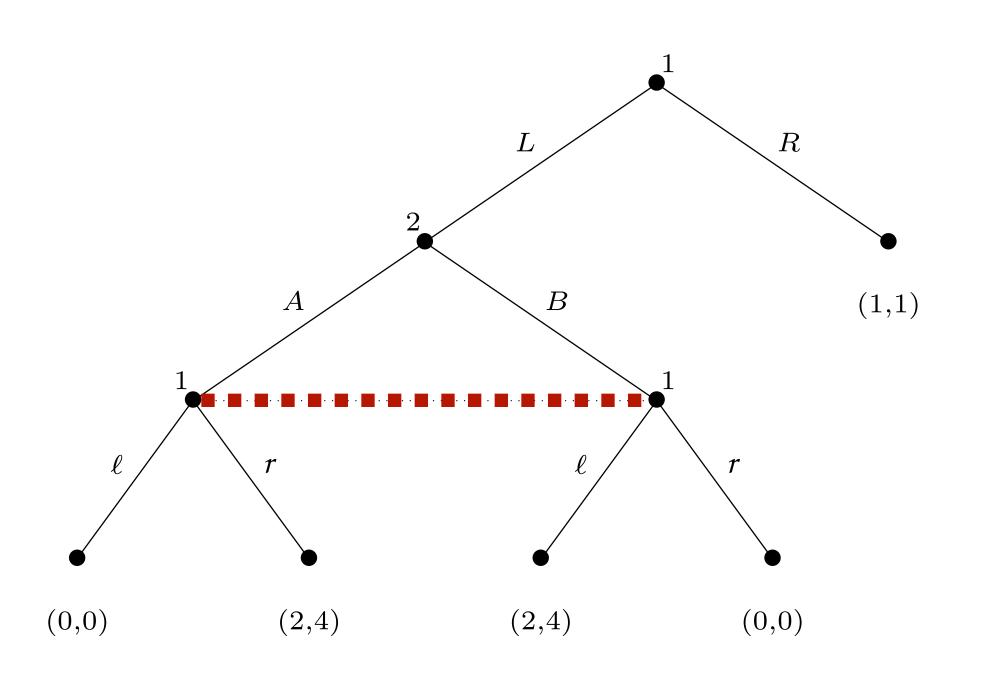
 A pure strategy associates an action with each information set, even those that will never be reached

Questions:

In an imperfect information game:

- What are the mixed strategies?
- 2. What is a best response?
- 3. What is a Nash equilibrium?

Induced Normal Form



1	Α	В
L,ℓ	0,0	2,4
L,r	2,4	0,0
R,ℓ	1,1	1,1
R,r	1,1	1,1

Question:

Can you represent an arbitrary perfect information extensive form game as an imperfect information game?

- Any pair of pure strategies uniquely identifies a terminal node, which identifies a utility for each agent
- We have now defined a set of agents, pure strategies, and utility functions
- Any extensive form game defines a corresponding induced normal form game

Summary

- Extensive form games model sequential actions
- Pure strategies for extensive form games map choice nodes to actions
 - Induced normal form: normal form game with these pure strategies
 - Notions of mixed strategy, best response, etc. translate directly
- Perfect information: Every agent sees all actions of the other players
 - Backward induction computes a pure strategy Nash equilibrium for any perfect information extensive form game
- Imperfect information: Some actions are hidden
 - Histories are partitioned into information sets; players cannot distinguish between histories in the same information set