Local Search

CMPUT 366: Intelligent Systems

P&M §4.7

Labs & Assignment #1

- Assignment #1 will be released today See eClass
- Due February 3 before lecture
- This week's lab: Thursday, 5:00pm to 8:00pm, <u>BS M 149</u> (this room!)
 - Not mandatory
 - You can get help from the TAs on your assignment in labs
 - Brief Python refresher at the beginning

Recap

- Graph search problems are an extremely general encoding for choosing a sequence of actions from a start state to a goal state
- Using heuristic functions can speed this process up
 - A* search is optimal but space-intensive
 - Branch & bound depth-first search is optimal and space efficient, but needs a good starting bound
- Varying the direction of search can exploit mismatches in forward and reverse branching factors

Lecture Outline

- 1. Recap & Logistics
- 2. Local Search
- 3. Hill Climbing
- 4. Randomized Algorithms

Searching for Goal Nodes

construct one.

Example (SAT problem): Given a Boolean formula,

 $P(X) = (X_1 \wedge X_2 \wedge \neg X_3)$

the formula true?

- **State** is the values of the different variables
- **Easy to recognize** when we've succeeded, but computing a "satisfying assignment" is **NP-complete** in general
- **SAT** is an example of a **constraint satisfaction problem**

- Sometimes, we know how to recognize a goal node, but not how to

$$\vee \ldots \vee (\neg X_{k-2} \land \neg X_{k-1} \land X_k),$$

is there an assignment of truth values to the variables X_i that makes

Searching for Goal Nodes

as states, variable value changes as actions), *but*:

- The **space is too big** to explore exhaustively
 - **Question:** How many states are there in a SAT problem with k variables?
 - Industrial SAT problems routinely have hundreds of thousands of variables
- 2. We don't care about the **sequence of actions**
 - Once we have a satisfying assignment, we are done

We can encode SAT as a graph search problem (assignments)

- Idea: start from a random assignment, and then search around in the space of possible assignments
- Need not keep track of the sequence of moves that we took
- Intuitively:
 - Select an assignment of a value to each variable
 - 2. Repeat:
 - Select a variable to change (İ)
 - Select a new value for that variable
 - 3. until a satisfying assignment is found

Local Search

Local Search Problem

Definition: Local Search Problem

- A constraint satisfaction problem: A set of variables, domains for the variables, and constraints on their joint assignment.
- Neighbours function: neighbours(n)
 - Maps from a node *n* to a set of "similar" nodes
- Score function: *score*(*n*)
 - Evaluates the "quality" of an assignment

Questions:

- 1. What are the nodes?
- 2. What are the goal nodes?



Neighbourhoods

- In previous graph search problems, the successor function represents states that can be reached from a given state by taking some actual action
 - In local search problems, the neighbours function is a design decision
 - We choose actions that will help us efficiently **explore the space** rather than trying to represent **actual actions**
- Usually the neighbourhood is states that differ in small ways from the current state (variable assignment)
 - E.g.: Assignments that differ in k different variables, possibly by a small amount
- Question: What might be a good neighbourhood function for SAT?

Heuristics vs. Scores

- Previously, the heuristic was optional, for improving efficiency
- In local search problems, the score function is required
 - The state space is **too big** to exhaustively explore, so uninformed search is not an option
 - Sometimes we don't even have a goal, we just want to maximize the quality of the state
- Example scores: number of unsatisfied clauses (in SAT); number of violated constraints (in CSP)

Generic Local Search Algorithm

Input: a constraint satisfaction problem; a *neighbours* function; a score function to maximize; a stop_walk criterion

current := random assignment of values to variables incumbent := current

repeat

if *incumbent* is a satisfying assignment:

return incumbent

if stop_walk():

else:

select a *current* from *neighbours*(*current*) **if** *score*(*current*) > score(incumbent):

incumbent := current until termination

- *current* := new random assignment of values to variables

Hill Climbing

- Idea: Select the neighbour with the highest score \bullet
 - This is called an **improving step**
 - If no improving steps available, halt and return *incumbent*
- We'll move toward the best solution once we are close \bullet enough
- This algorithm is called **hill climbing**:
 - It seeks the highest point on the scoring function's graph
 - It moves only **uphill** (i.e., it makes only improving steps)

Hill Climbing Algorithm

Input: a constraint satisfaction problem; a *neighbours* function; a *score* function

current := random assignment of values to variables *incumbent* := *current*

repeat

if *incumbent* is a satisfying assignment: return incumbent

if *False*:

current := new random assignment of values to variables

else:

current := *n* from *neighbours*(*current*) with maximum *score*(*n*)

if score(current) > score(incumbent): *incumbent* := *current*

else:

return incumbent

until termination

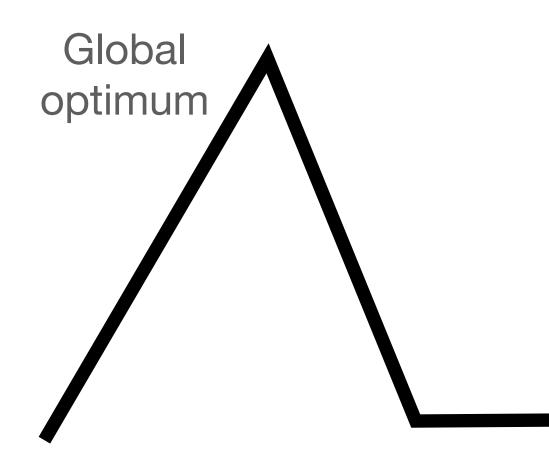
Questions:

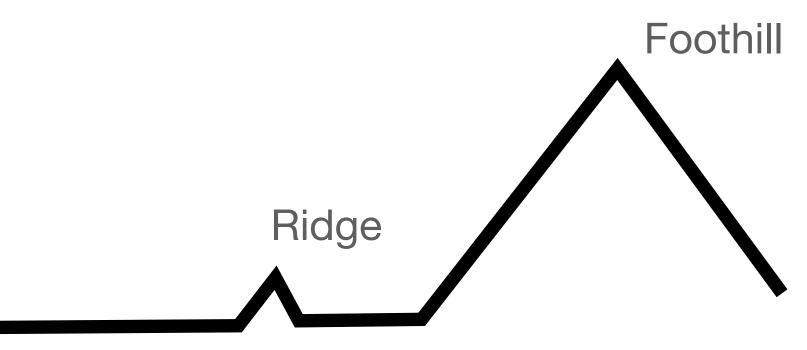
- 1. Is hill climbing complete?
- 2. Is hill climbing optimal?



Hill Climbing Problems

- 1. Foothills: Local maxima that are not global maxima
- 2. Plateaus: Regions of the state space where the score is uninformative
- 3. Ridges: Foothills that would not be foothills with a larger neighbourhood
- 4. **Ignorance of the global optimum:** Unless we reach a satisfying assignment, we cannot be sure that an optimum returned by local search is the **global optimum**.





Randomized Algorithms

- Adding random moves can fix some hill climbing problems
- Two main kinds of random move:
 - 1. **Random restart:** Start searching from a **completely** random new location
 - 2. Random step: Choose a random neighbour
- Stochastic random search: Add both kinds of random moves to hill climbing

Input: a constraint satisfaction problem; a *neighbours* function; a *score* function to maximize; a stop_walk criterion; a random_step criterion

current := random assignment of values to variables *incumbent* := *current*

repeat

if *incumbent* is a satisfying assignment: return incumbent

if stop_walk():

current := new random assignment of values to variables else if random_step():

current := a random element from *neighbours*(*current*)

else:

current := n from neighbours(current) with maximum score(n) **if** *score*(*current*) > score(incumbent): incumbent := current

Stochastic Local Search

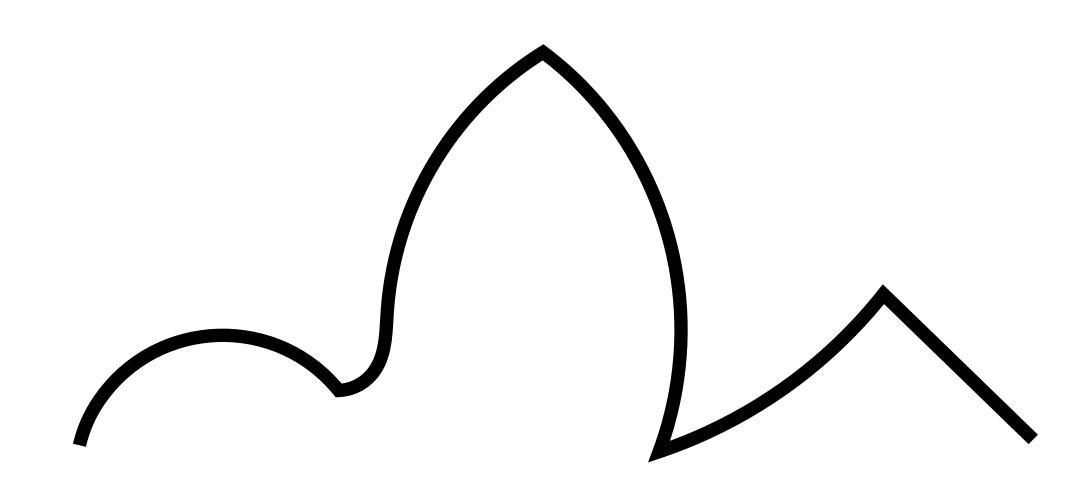
Questions:

- Is stochastic local search complete? (**Why?**)
- 2. Is stochastic local search optimal? (Why?)



Two Examples

- Consider two partial algorithms:
 - 1. Hill climbing plus random restart
 - 2. Hill climbing plus random steps
- **Question:** Which finds the maximum most easily on each of these two search spaces? Why?



Simulated Annealing

- Idea: Start out by searching pretty randomly, but become more directed
 - **Intuition:** Move to a good neighbourhood quickly, then search intensively \bullet in that neighbourhood
- Maintain a "temperature" T lacksquare
- Choose new nodes more randomly at higher temperatures;
- At each step:
 - Randomly choose a neighbour *new*

Gradually decrease the temperature (according to a **cooling schedule**)

2. Always accept (i.e., assign to *current*) if score(new) > score(current)3. Else, accept with **probability** $e^{[(score(new)-score(current))/T]}$

Summary

- For some problems, we only care about finding a **goal node**, not the actions we took to find it
- Local search: Look for goal states by iteratively moving from a current state to a neighbouring state
 - Hill climbing: Always move to the highest-score neighbour
 - Random step: Sometimes choose a random neighbour
 - Random restart: Sometimes start again from an entirely random state
 - Simulated annealing: Random moves start very random, become more greedy over time